

Non-precise def.: 1-dim geom object described by polynomial equations.

Def. (for this lecture): Algebraic curve:  $C \subset \mathbb{C}^2$ ,  $C = \{(x,y) \in \mathbb{C}^2 \mid f(x,y) = 0\}$ ,  $f \in \mathbb{C}[x,y]$  non-constant.

Examples.

$f$	$C \cap \mathbb{R}^2$	$\deg(f)$
$x+y+1$		1
$x^2 + y^2 - 1$		2
$(y^2 - x^3)^c - 4x^5y + x^6 - x^7$		7
$y^2 - x^3 + x$		3
$x^2 + y^2 + 1$		2

Def.: Let  $p \in \mathbb{C}[x_1, \dots, x_n]$ , then  $\boxed{\deg p} := \min \{n \in \mathbb{N} \cup \{-\infty\} \mid p = \sum a_{i_1, \dots, i_k} x_1^{i_1} \dots x_k^{i_k} \text{ & } i_1, \dots, i_k \leq n\}$

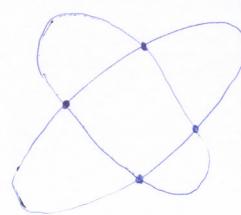
How do we see the degree geometrically?  $\rightarrow$  intersect with a (general) line, and count the intersection points

More generally:  $C, C' \subset \mathbb{C}^2$  alg curves def'd by  $f, g \in \mathbb{C}[x,y]$ ,  $\deg f = m$ ,  $\deg g = n$

Then what can be said about  $C \cap C' \subset \mathbb{C}^2$ ? Specifically,  $\#(C \cap C') = ?$

Ex.  $f = y - x^3$   
 $g = y - x$

$C'$  is a curve



2 ellipses

$$\#(C \cap C') = 3 = 3 \cdot 1$$

$$\#(C \cap C') = 4 = 2 \cdot 2$$

$$\#(C \cap C') = 6 = 3 \cdot 2$$

But 2 circles may have 2 intersections:



Or even 0:



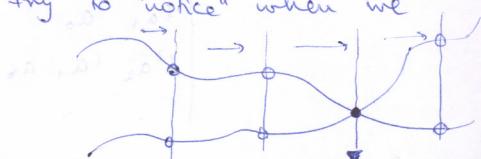
Still, it seems like  $\#(C \cap C') \leq \deg f \cdot \deg g$

Pathological case:  $C = C' \Rightarrow \#(C \cap C') = \infty$

Theorem (Bézout):  $C, C' \subset \mathbb{C}^2$  alg curves as above. If  $(f, g) = 1$  then  $\#(C \cap C') \leq m \cdot n$ .

PF: Idea: scan through the curves with a line, and try to "notice" when we hit an intersection point.

$$f(x,y) = a_0(x)y^m + \dots + a_m(x), \quad g(x,y) = b_0(x)y^n + \dots + b_n(x)$$



When do  $f(x_0, y)$  and  $g(x_0, y) \in \mathbb{C}[y]$  have a common root?

(This classmate's notes mention that this is related to the gcd of the polynomials.)

Exercise.  $X := \{(x_1, y_1, z) \in \mathbb{C}^3 \mid xy = yz = zx = 0\}$  Show that  $\nexists$   $f, g \in \mathbb{C}[x, y, z]$  s.t.  $f(x_1, y_1, z_1) = 0$  &  $g(x_1, y_1, z_1) = 0$

Hint:  $\mathbb{P}^2 \mathbb{C}$  + Bézout.

For  $F, G \in \mathbb{C}[x]$  if  $F(a) = G(a) = 0$  for some  $a \in \mathbb{C}$  then  $F(x) = (x-a) \cdot u(x)$   $G(x) = (x-a) \cdot v(x)$

$\Rightarrow vF - uG = 0$  and  $\deg u < \deg F$ ,  $\deg v < \deg G$

The converse holds too.  $\rightarrow$  We can use this to detect common roots of  $f$  and  $g$ .

$\rightarrow$  Solve  $vF - uG = 0$  for  $u, v$  where

$$u(x) = c_0 x^{m-1} + \dots + c_{m-1}$$

$$v(x) = d_0 x^{n-1} + \dots + d_{n-1}$$

given to be found

$$vF - uG = (a_0 d_0 - b_0 c_0) \cdot x^{m+n-1} + (a_1 d_0 + a_0 d_1 - b_1 c_0 - b_0 c_1) x^{m+n-2} + \dots$$

$\rightarrow$  yields  $m+n$  linear equations in the variables  $c_0, \dots, c_{m-1}, d_0, \dots, d_{n-1}$

$$\begin{matrix} n \text{ columns} & m \text{ columns} \\ \begin{bmatrix} a_0 & | & b_0 \\ a_1 & a_0 & | & b_1 & | & \dots & b_n \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_m & | & b_m \\ a_m & | & b_n \\ \vdots & \vdots & \vdots \\ a_m & | & b_n \end{bmatrix} & \begin{bmatrix} d_0 \\ \vdots \\ d_{n-1} \\ c_0 \\ \vdots \\ c_m \end{bmatrix} \end{matrix} = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

The det of this mtx is the resultant  $R_{F,G}$

because  $(f, g) = 1$

If  $R_{F,G}(x) = 0$  then  $\exists y \in \mathbb{C} : f(x, y) = g(x, y)$ . Since  $R_{F,G}$  is a non-constant polynomial in  $x$ , there are only fin many  $(x_1, \dots, x_k)$  s.t.  $f(x_i, y) = g(x_i, y)$ .  $\Rightarrow \#(\mathcal{C} \cap \mathcal{C}') < \infty$

By rotating wma there is at most one intersection point on each vertical line.

Now we estimate  $\deg R_{F,G}$ , which will give us an upper bound for  $\#(\mathcal{C} \cap \mathcal{C}')$ .

Claim.  $\deg R_{F,G} \leq mn$

$$\text{PF: } R_{F,G}(tx) = \begin{bmatrix} a_0 & & & & \\ ta_1 & a_0 & & & \\ & \ddots & \ddots & \ddots & \\ & & t^2 a_2 & ta_1 & a_0 \\ & & & \vdots & \vdots \end{bmatrix} = t^{m+n} \cdot R_{F,G}(x) \text{ whenever } a_i(x) = x^i, b_j(x) = x^j$$

Notice that  $\deg a_i \leq i$ ,  $\deg b_j \leq j$ . Due to degree counting:  $\deg a_i \leq i$

$$(\deg a_i y^{m-i} = \deg a_i + \deg y^{m-i} \leq m \Rightarrow \deg a_i(x) \leq i)$$

We get the exponent  $m \cdot n$  as  $(1 + \dots + (m+n-1)) - (1 + \dots + (m-1)) - (1 + \dots + (n-1))$   
(matrix operations)

For general  $a_i, b_j \rightarrow$  only the highest degree matters.

Consider  $p_1(x), \dots, p_k(x) \in \mathbb{R}[x]$ , and draw their graphs.

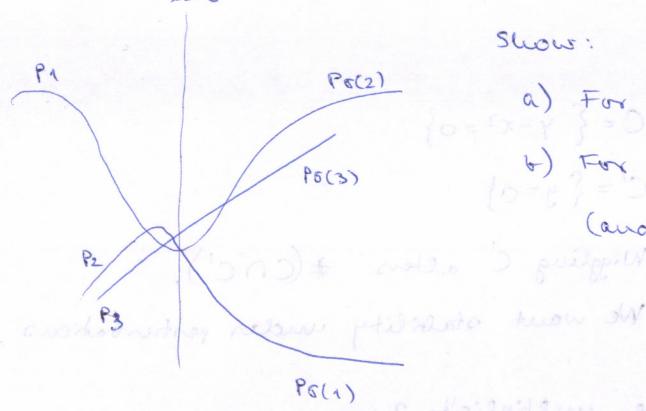
Exercise.<sup>4</sup> For  $p_1, \dots, p_k$  pairwise distinct, define a permutation:

For small  $x < 0$ ,  $p_1(x) < \dots < p_k(x)$  may be assumed.

Then for small  $x > 0$ ,  $p_{\sigma(1)}(x) < \dots < p_{\sigma(k)}(x)$  for some  $\sigma \in S_k$

In this context, small means small in absolute value; more precisely:

$$\exists a > 0 \text{ s.t. } |x| < a \text{ and } \forall y \in \mathbb{R}, |y| < a : p_i(y) \neq p_j(y) \quad \forall i \neq j.$$



Show :

- a) For  $k \leq 3$ , every permutation can be realised.  
 b) For  $k = 4$ , there is one that cannot be realised  
 (and find such a permutation).

Exercise.  $p \in \mathbb{R}[x]$ ,  $p(x) > 0 \quad \forall x \in \mathbb{R} \Rightarrow 2|\deg p| \leq p$  has a global minimum. (known)

To generalise, consider  $p(x_1y) = x^2y^2 + 2xy + x^2 + 1 \in \mathbb{R}[x_1y]$ . Show  $p(x_1y) > 0$

$f(x,y) \in \mathbb{R}^2$  but has no minimum.

Exercise 5b. *monotonic & local max*  
 Reminder: if a polynomial in 1 variable has 2 local minima then  
 it has a local maximum too.

Show that  $p(x,y) = (x^2y - x - 1)^2 + (x^2 - 1)^2$  has 2 global minima, but no other critical points.

Exercise? Where are the missing pts?

Recall Exercise 4 from last time.

Hint: This cannot be realised:  $(1, 2, 3, 4) \mapsto (2, 4, 1, 3)$ .

Any  $f \in \mathbb{R}[x]$  can be written as  $f(x) = a_m x^m + \dots + a_1 x + a_0$ , for  $a_m, a_1 \neq 0$ .

$$\begin{aligned} m_0(f) &:= m && \text{multiplicity of } f \text{ at } 0 / \text{valuation of } f \\ \deg(f) &:= d && \text{degree of } f \end{aligned}$$

Note that for  $|x| \gg 0$ ,  $f(x) \sim a_d x^d$ ,

for  $|x| \ll 1$ ,  $f(x) \sim a_m x^m$ .

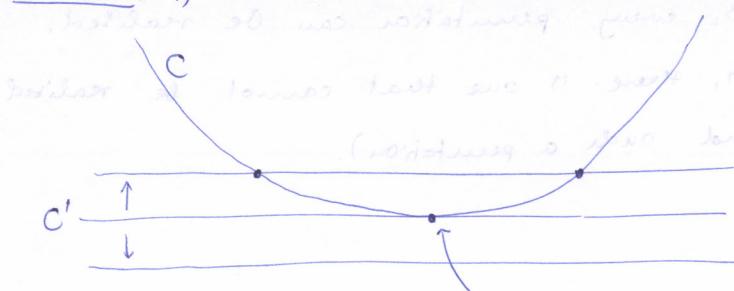
If  $2|m_0(f)|$ ,  $f$  stays on the same side of the  $x$ -axis,

if  $2|m_0(f)|$ ,  $f$  changes sides.

If  $m_0(f) > m_0(g)$  then for  $|x| \ll 1$ ,  $|f(x)| < |g(x)|$

Goal: Get equality in Bézout.

Problems: 1)



$$C = \{y - x^2 = 0\}$$

$$C' = \{y = 0\}$$

Wiggling  $C'$  alters  $\#(C \cap C')$ .

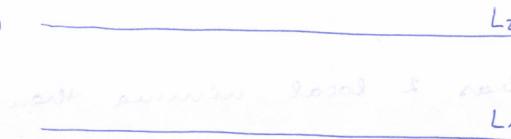
We want stability under perturbations.

Want this to have multiplicity 2:

$$\begin{cases} y - x^2 = 0 \\ y = 0 \end{cases} \Rightarrow x^2 = 0$$

Solution: define intersection multiplicities.

2)



Should have 1 intersection.

→ we wish for points at infinity

→ projective space.

Projective spaces

$$\mathbb{P}^n(\mathbb{C}) := \mathbb{C}^{n+1} \setminus \{0\} / \mathbb{C}^\times$$

$$\mathbb{P}^n(\mathbb{R}) := \mathbb{R}^{n+1} \setminus \{0\} / \mathbb{R}$$

Ex.  $\mathbb{P}^1(\mathbb{R}) \cong S^1$ ,  $\mathbb{P}^3(\mathbb{R}) \cong SO(3)$  as groups,  $\mathbb{P}^1(\mathbb{C}) \cong S^2$  diffeomorphism  
(exercise)

$$\varphi_i: \mathbb{C}^n \hookrightarrow \mathbb{P}^n(\mathbb{C})$$

$$(x_1, \dots, x_n) \mapsto [x_1 : \dots : x_{n-1} : 1 : x_n : \dots : x_n]$$

$$\mathbb{P}^n(\mathbb{C}) \setminus \varphi_i(\mathbb{C}^n) = \left\{ [x_0 : \dots : 0 : \dots : x_n] \mid x_0, \dots, x_n \in \mathbb{C} \text{ not all } 0 \right\} \cong \mathbb{P}^{n-1}(\mathbb{C})$$

Def: Projective algebraic curve:  $C \subset \mathbb{P}^2(\mathbb{C})$ ,  $C = \{(x:y:z) \in \mathbb{P}^2(\mathbb{C}) \mid F(x,y,z) = 0\}$

where  $F \in \mathbb{C}[x,y,z]$  is nonconstant & homogeneous.

Homogeneous polynomials, homogenisation.

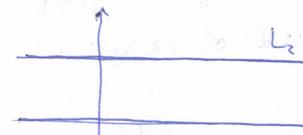
$$\{C \subset \mathbb{C}^2\} \longleftrightarrow \{\bar{C} \subset \mathbb{P}^2(\mathbb{C})\}$$

$$C = \{f = 0\} \longleftrightarrow \{F = 0\} = \bar{C} \quad \text{where } F \text{ is the homogenisation}$$

Ex:  $L_1 = \{(x,0)\}, L_2 = \{(x,1)\} \subset \mathbb{C}^2$

$$f_1(x,y) = y \quad f_2(x,y) = y - 1$$

↓ homogenise



$$F_1(x,y,z) = y$$

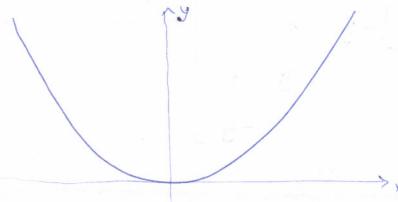
$$F_2(x,y,z) = y - z$$

→ intersections in  $\mathbb{P}^2(\mathbb{C})$ :

$$\begin{cases} y=0 \\ y-z=0 \end{cases} \Rightarrow y=z=0 \Rightarrow x \neq 0 \Rightarrow [1:0:0]$$

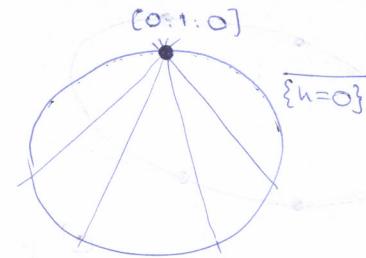
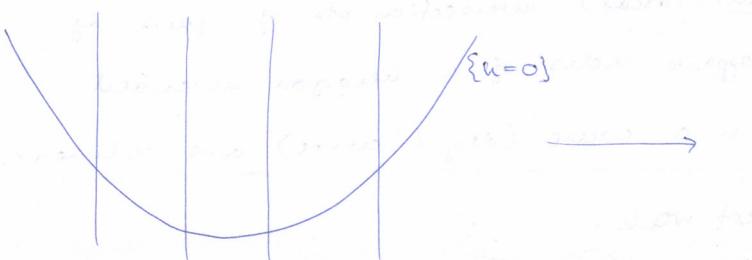
the point is in  
the direction of the  
x-axis

Ex:  $h(x,y) = y - x^2 \rightarrow yz - x^2$



point(s) at infinity:

$$[0:1:0] \in \overline{\{h=0\}} \subset \mathbb{P}^2(\mathbb{C})$$



Theorem. (Bézout in  $\mathbb{P}^2$ )  $C, C' \subset \mathbb{P}^2(\mathbb{C})$  algebraic curves without common components.

$$\Rightarrow \#_{\text{mult}}(C \cap C') := \sum_{P \in C \cap C'} i_p(C, C') = \deg C \cdot \deg C'$$

where  $i_p(C, C') \in \mathbb{Z}$  is the intersection multiplicity.

Do the same as before:  $F(x,y,z) = \sum_{i=0}^m a_i(x,y)z^{m-i}$   $a_i, b_i \in \mathbb{C}[x,y]$  homogeneous of deg  $i$

$$G(x,y,z) = \sum_{i=0}^n b_i(x,y)z^{n-i}$$

$$\det \begin{bmatrix} a_0(x,y) & 0 & -b_0(x,y) & 0 \\ \vdots & \vdots & \vdots & \vdots \\ a_m(x,y) & a_0(x,y) & -b_m(x,y) & -b_0(x,y) \\ \vdots & \vdots & \vdots & \vdots \\ 0 & a_m(x,y) & -b_n(x,y) & 0 \end{bmatrix} =: R_{F,G}(x,y) \in \mathbb{C}[x,y]$$

homog of deg m+n

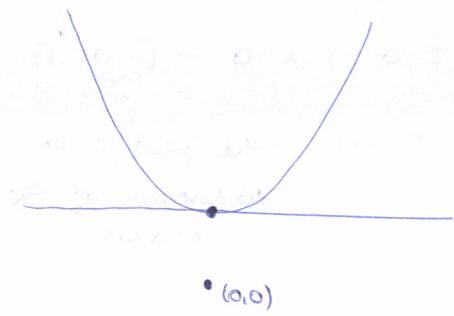
Def. For  $P \in C \cap C'$ , let  $i_p(C, C')$  be the multiplicity of the zero  $[x:y]$  of  $R_{F,G}(x,y)$ .

Alternatively:  $i_p(C, C') = \begin{cases} \text{mult. of } [1:y/x] \text{ in } R_{F,G}(1,y) & \text{if } x \neq 0, \\ \text{mult. of } [x/y:1] \text{ in } R_{F,G}(x,1) & \text{if } y \neq 0 \end{cases}$

Ex.  $C = \{(x:y:z) \in \mathbb{P}^2(\mathbb{C}) \mid yz - x^2 - z^2 = 0\}$

$$C' = \{(x:y:z) \in \mathbb{P}^2(\mathbb{C}) \mid y - z = 0\}$$

$$C \cap C' = \{P = [0:1:1]\}$$



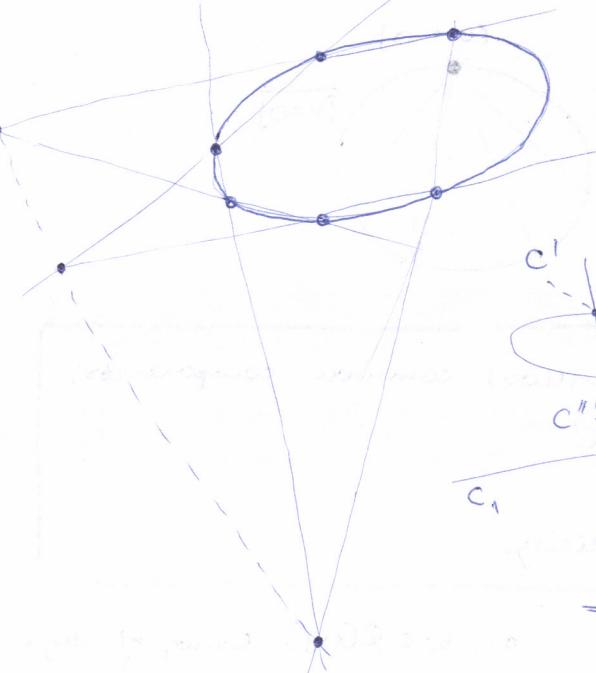
$$F(x,y,z) = (-1)z^2 + y \cdot z^1 - x^2 \cdot z^0$$

$$G(x,y,z) = z^1 - y \cdot z^0$$

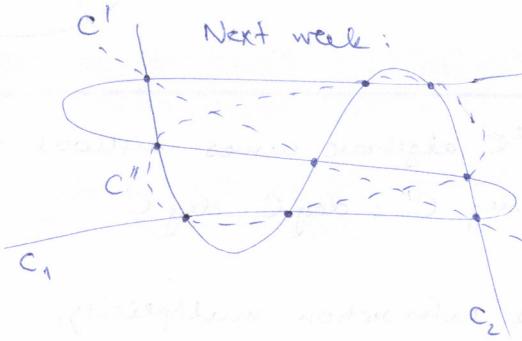
$$R_{F,G}(x,y) = \det \begin{bmatrix} -1 & 1 & 0 \\ y & -y & 1 \\ -x^2 & 0 & -y \end{bmatrix} = -x^2$$

$$\Rightarrow i_p(C, C') = 2$$

Thm. (Pascal) Intersection pts of pairs of opposite sides of a hexagon inscribed in a conic (deg 2 curve) are collinear.



Next week:

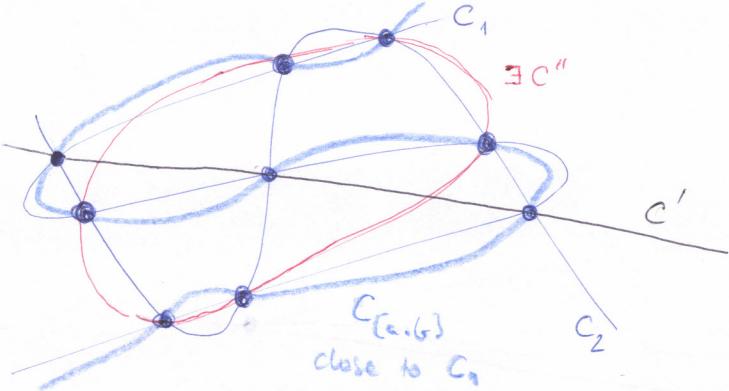


$$\deg C_1 = \deg C_2 = d$$

$$C_1 \cap C_2 = \{d^2 \text{ points}\}$$

$\deg C' = k$  irreducible,  
and  $C''$  contains  $k \cdot d$  of these pts

$$\Rightarrow \exists C'', \deg C'' = d-k \text{ going through the rest.}$$



$$\deg C_1 = \deg C_2 = d$$

$$\deg C' = k \quad \text{irreducible}$$

$$\deg C'' = d-k$$

Lemma.  $C_1, C_2 \subset \mathbb{P}^2(\mathbb{C})$  curves of degree  $d$ ,  $C_1 \cap C_2 = \{d^2 \text{ points}\}$ .

Assume there is an irreducible curve  $C'$  of degree  $k < d$ , passing through  $k \cdot d$  pts of  $C_1 \cap C_2$ . Then there exists a curve  $C''$  of degree  $d-k$  through the remaining  $(d-k) \cdot d$  points.

PF:  $C_i = \{F_i = 0\}$ ,  $F_i \in \mathbb{C}[x,y,z]$  homogeneous of deg  $d$

For  $[a:b] \in \mathbb{P}^1(\mathbb{C})$  consider the curve  $C_{[a:b]} := \{aF_1 + bF_2 = 0\}$

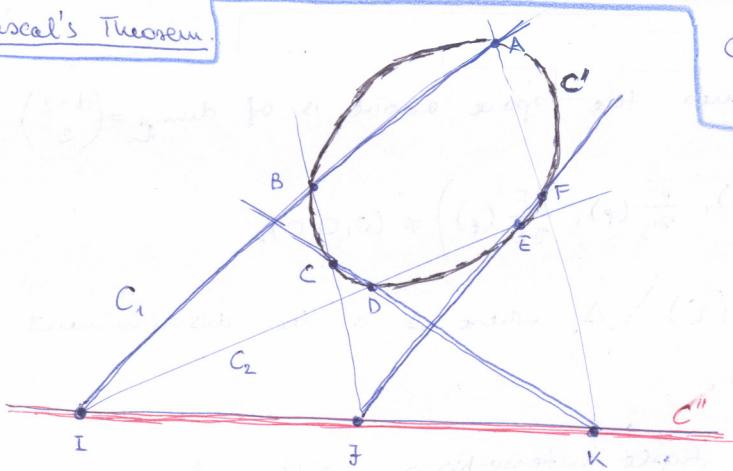
Then  $\deg C_{[a:b]} = d$  and  $\forall [a:b] \in \mathbb{P}^1(\mathbb{C}): (C_1 \cap C_2) \subset C_{[a:b]}$ . Think of this as an algebraic version of an isotopy between curves.

Let  $p \in C' \setminus (C_1 \cap C_2)$ ,  $a := F_2(p)$ ,  $b := F_1(p) \Rightarrow [a:b] \in \mathbb{P}^1(\mathbb{C})$ ,

$C' \cap C_{[a:b]}$  contains at least  $k \cdot d + 1$  points  $\Rightarrow C'$  and  $C_{[a:b]}$  share a common component by Bézout.

Since  $C'$  is irreducible:  $C_{[a:b]} = C' \cup C''$  for some curve  $C''$  of degree  $d-k$ .

Pascal's Theorem.



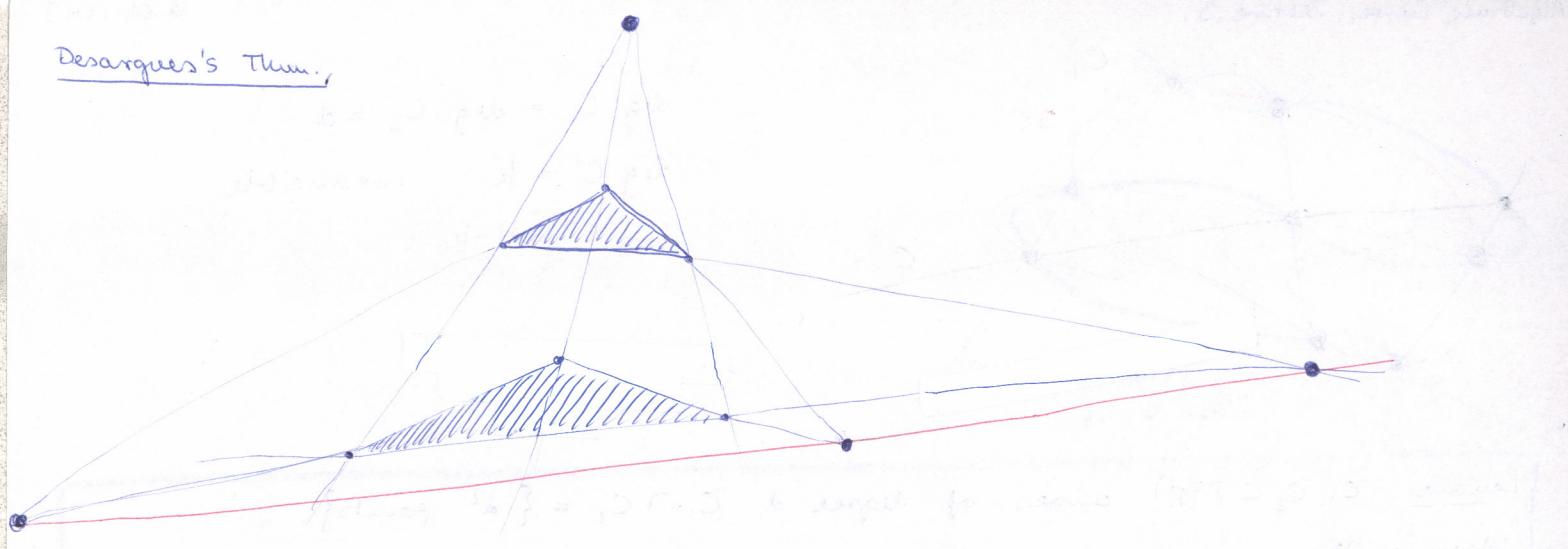
Claim:  $I, J, K$  are collinear, where  $C'$  is a smooth conic in  $\mathbb{P}^2(\mathbb{C})$

PF: Use the Lemma with

$$C_1 := \overline{AB} \cup \overline{CD} \cup \overline{EF}$$

$$C_2 := \overline{BC} \cup \overline{DE} \cup \overline{FA}$$

Desargues's Thm.



Go one dimension higher, look at the planes of the triangles.

### Smooth curves

Recall: connected compact oriented surfaces are classified by their genera.



$$g=0 \quad g=1 \quad g=2$$

The genus tells us how many "broken glasses" we can embed (This is how the classification result is proven.)

Fix some  $d \in \mathbb{N}$  and consider  $\mathcal{C}_d := \{ C = \{ F(x,y,z) = 0 \} \subset \mathbb{P}^2(\mathbb{C}) \mid F \in \mathbb{C}[x,y,z] \text{ homog of deg } d \}$

$$F(x,y,z) = \sum_{i+j+k=d} a_{ijk} x^i y^j z^k \quad \text{for some } a_{ijk} \in \mathbb{C} \text{ not all zero.}$$

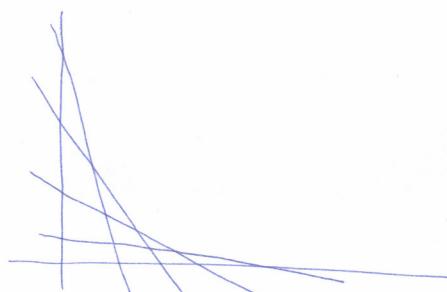
There are  $\binom{d+2}{2}$  terms in this  $\sum$ , thus the space above is of dim<sub>C</sub> =  $\binom{d+2}{2}$ .

Def.  $C = \{ F = 0 \}$  is smooth if  $\forall p \in C: \left( \frac{\partial F}{\partial x}(p), \frac{\partial F}{\partial y}(p), \frac{\partial F}{\partial z}(p) \right) \neq (0,0,0)$ .

$\{ \text{smooth curves of degree } d \} \cong \mathbb{P}^{\binom{d+2}{2}-1}(\mathbb{C}) \setminus \Delta$  where  $\Delta$  is the discriminant.

Note that this space is path-connected.

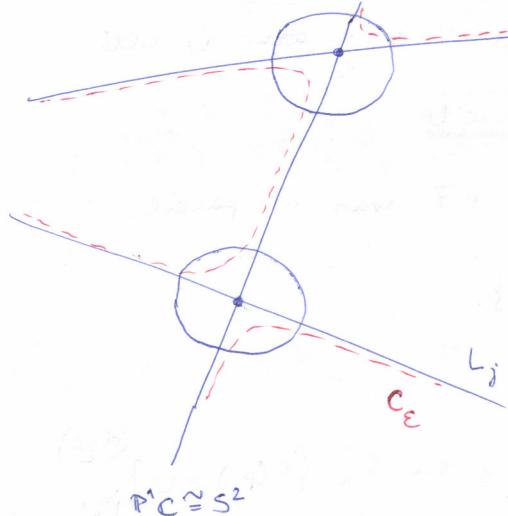
Let  $L_1, \dots, L_d \subset \mathbb{P}^2(\mathbb{C})$  be lines without triple intersections, given by  $l_1, \dots, l_d \in \mathbb{C}[x,y]$  linear forms.



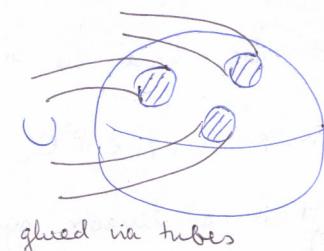
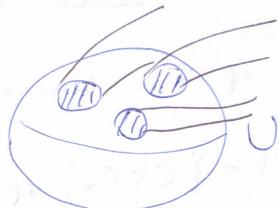
$$\text{Consider } C_0 := \{ l_1 \cdot l_2 \cdot \dots \cdot l_d = 0 \}$$

and perturb it slightly:  $C_\epsilon := \{ l_1 \cdots l_d - \epsilon l_0^d = 0 \}$  where  $l_0$  is chosen s.t.  $C_\epsilon$  becomes smooth for small  $\epsilon > 0$ .

Such an  $\ell_\varepsilon$  exists: it suffices for it not to vanish at any of the intersections.

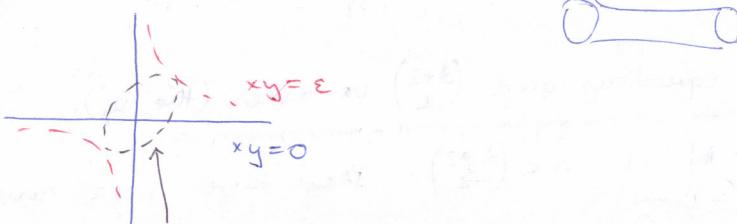


This is  $S^2 \setminus (d-1 \text{ disks})$  d times:



glued via tubes

$\{xy = \varepsilon\} \cap B(O, R) \cong \text{cylinder}$



this circle is visible only over  $C$ ,  
red and black together give a cylinder

$$\begin{aligned}\chi(C_\varepsilon) &= d \cdot \chi(S^2 \setminus (d-1 \text{ disks})) + \dots + \underbrace{\chi(\text{tube})}_{0} - \dots - \underbrace{\chi(S^1)}_{0} \\ &= d \cdot (2 - d + 1) = -d^2 + 3d\end{aligned}$$

$$g(C_\varepsilon) = \frac{(d-1)(d-2)}{2}$$

Notice that these depend on  $d = \deg F$  only.

Recall the correspondence  $C_d \longleftrightarrow \mathbb{P}^{(d+2)-1} \mathbb{C}$

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$$C = \{F=0\} \longleftrightarrow [a_0 : \dots : a_{\binom{d+2}{2}}]$$

where  $F = \sum_{i=1}^{\binom{d+2}{2}} a_i M_i$  where  $M_i$  is a monomial

Notice that this is rather unrigorous: e.g. groups don't form a group, but here we have that curves form a projective space.

Def: A **linear system of degree  $d$  curves** is a subset  $V \subseteq C_d$  whose image in  $\mathbb{P}^{(d+2)-1} \mathbb{C}$  is given by linear equations.

Special cases:  $C_1 = \{\text{lines}\} \longleftrightarrow \mathbb{P}^2 \mathbb{C}$  this is how  $\mathbb{P}^2$  was def'd

$$C_2 = \{\text{quadratics}\} \longleftrightarrow \mathbb{P}^5 \mathbb{C}$$

$$C_3 = \{\text{cubics}\} \longleftrightarrow \mathbb{P}^9 \mathbb{C}$$

Example. Fix a conic  $C \subset \mathbb{P}^2$ , and let  $V := \{T_p C \mid p \in C\}$  be the set of tangent lines.  $V \subset \mathbb{C}_1 \cong \mathbb{P}^1$ . Ex.: show that  $V$  is a conic as well.

Def. A linear system of dimension 1 is called a **pencil**.

E.g. what we had in the proof of the Lemma on p.7 was a pencil.

Example. Fix  $p_1, \dots, p_n \in \mathbb{P}^2$ ,  $d \in \mathbb{N}$ ,  $H := \{C \in \mathcal{C}_d \mid \forall p_i \in C\}$ .

Then  $H$  is a linear system.

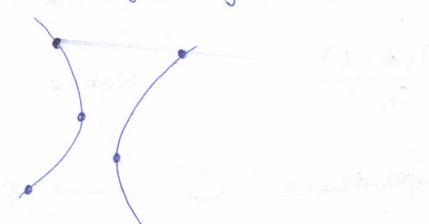
$F = \sum_{j=1}^{\binom{d+2}{2}} a_j M_j$ , we get a system of equations given by  $\{F(p_j) = 0\}_{j=1}^{\binom{d+2}{2}}$

We have  $n$  equations and  $\binom{d+2}{2}$  variables (the  $a_j$ ).

Lemma. Let  $n, d \in \mathbb{N}$  s.t.  $n < \binom{d+2}{2}$ . Then there is a curve of degree  $d$  through any set of  $n$  points. (Just linear algebra.) □

Ex.  $2 < \binom{1+2}{2}$   $\rightarrow$  there is a line through any 2 points

$5 < \binom{2+2}{2}$   $\rightarrow$  there is a quadric through any 5 points



We return to discussing smooth curves.

Recall  $\{C \in \mathcal{C}_d \mid C \text{ smooth}\} \hookrightarrow \mathbb{P}^{\binom{d+2}{2}-1}$

where  $\Delta$  is a hypersurface called the discriminant.

This is a purely theoretical result, in practice it is not at all convenient to do computations with  $\Delta$ . The important thing is that  $\mathbb{P}^{\binom{d+2}{2}-1} \setminus \Delta$  is path-connected.

Thm.  $C \in \mathcal{C}_d$  smooth  $\Rightarrow$   $C$  is a compact Riemann surface of genus

$$g(C) = \frac{(d-1)(d-2)}{2}$$

<u>Rmk.</u> There are gaps:	d	1	2	3	4	5	...
	g	0	0	1	3	6	...

In particular, there are no curves of genus 2.

There is a converse to this Thm:

Thm. (Chow, 1949)  $S \subset \mathbb{P}^2(\mathbb{C})$  compact Riemann surface (i.e. locally given by analytic equations). Then  $S$  is an algebraic curve (i.e. globally given by a polynomial equation).

This is truly remarkable: we get that  $S$  itself is an algebraic curve, not just isotopic (or 3rd like that) to one.

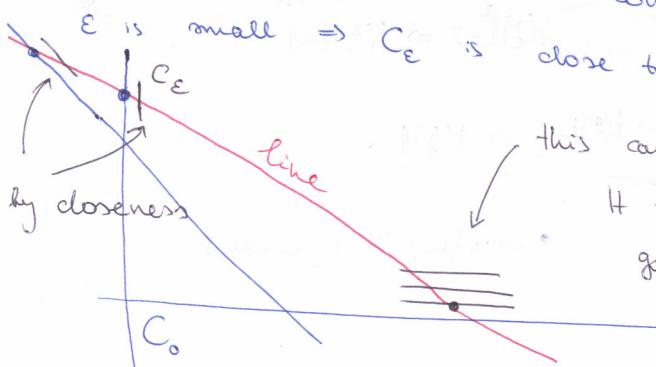
Ex.  $C := \{y = p(x)^2\} \subseteq \mathbb{C}^2$ ,  $p(x) = (x - a_1) \cdots (x - a_{2g+1})$ ,  $a_i \in \mathbb{C}$  distinct  
 $\Rightarrow C$  is a non-smooth surface of genus  $g$ .

In particular, we can realize all the genera. Why is this not in contradiction with the Rmk. above? (What happens when we projectivize?)

Thm. (degree-genus formula)  $C \in \mathbb{C}_d$  smooth  $\Rightarrow g(C) = \frac{(d-1)(d-2)}{2}$

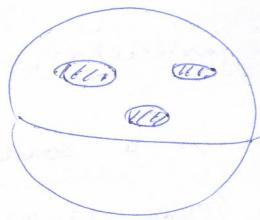
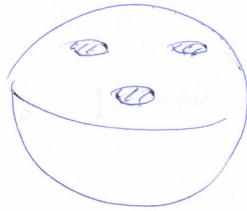
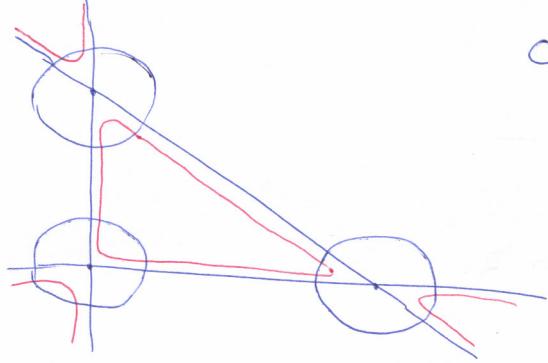
PF SKETCH: Consider  $C_0: l_1 l_2 \cdots l_d = 0$  where  $l_i$  are generic linear forms.  
 $L_i = \{l_i = 0\}$   $C_\epsilon: l_1 l_2 \cdots l_d = \epsilon \cdot l_0$  for  $\epsilon > 0$  small s.t.  $C_\epsilon$  is smooth of degree  $d$ .

Idea: Choose a path in  $\mathbb{P}^{(d+2)-1} \setminus \Delta$  connecting  $C$  and  $C_\epsilon \Rightarrow g(C) = g(C_\epsilon)$ .  
this cannot happen: there can't be multiple layers.  
It would contradict Bezout: have a line go through them and count the intersections.



What happens near intersection points?

$$C_\varepsilon \setminus \bigcup_{p \in L_i \cap L_j} B(p, R) = \coprod_{p \in L_i \cap L_j} (\mathbb{P}^1 \setminus \{(d-1)\text{ discs}\})$$



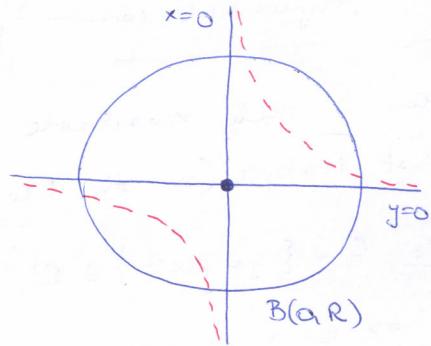
$d$  spheres with  $d-1$  discs removed from each  
Near an intersection point  $p \in L_i \cap L_j$  we may choose coordinates

$$\mathbb{C}^2 \hookrightarrow \mathbb{P}^2 \mathbb{C} \text{ s.t. } C_\varepsilon \cap B(p, R) \text{ is described by}$$

$$\{(x, y) \in \mathbb{C}^2 \mid xy = \varepsilon, |x|^2 + |y|^2 \leq R\}$$

$$4xy = (x+y)^2 - (x-y)^2$$

Change of coordinates:  $u := x+y, v := x-y,$



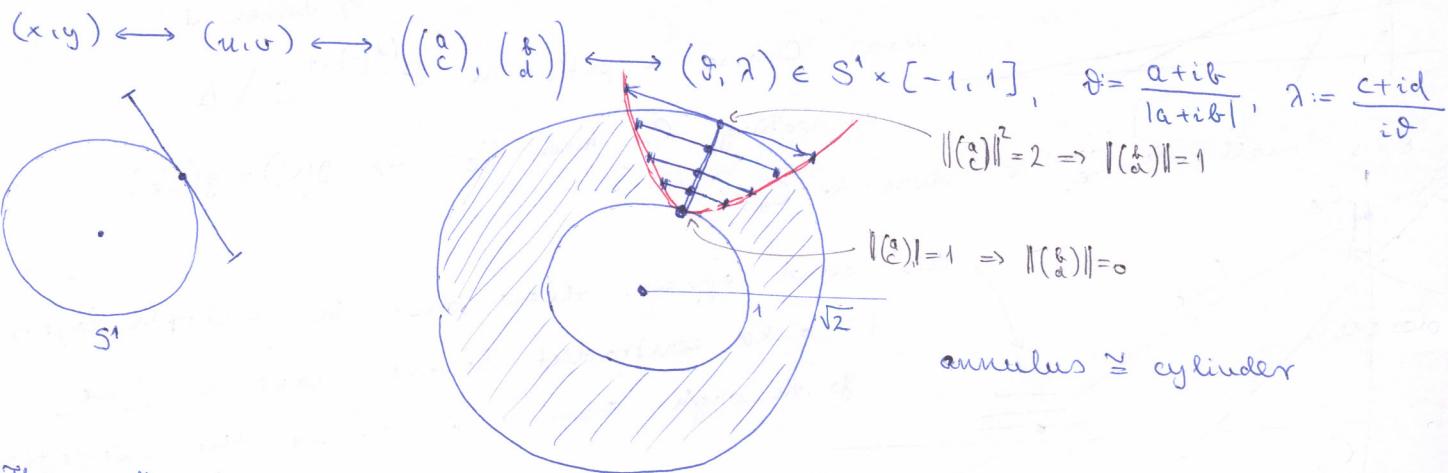
$xy = \varepsilon$  becomes  $u^2 + v^2 = 4\varepsilon$ . Write  $u = a+ib, v = c+id$ ,  $a, b, c, d \in \mathbb{R}$

$$\underbrace{a^2 + c^2 - b^2 - d^2}_{\|(\frac{a}{c})\|^2} + 2i(ab + cd) = 4\varepsilon$$

$$\|(\frac{a}{c})\|^2 - \|(\frac{b}{d})\|^2 = 0 \text{ as } (\frac{a}{c}) \perp (\frac{b}{d})$$

Rescale coordinates:  $2 \geq \|(\frac{a}{c})\|^2 = \|(\frac{b}{d})\|^2 + 1 \geq 1$

$\Rightarrow (\frac{a}{c}) \neq (\frac{b}{d})$  and  $\|(\frac{b}{d})\|^2 \in [0, 1]$ .  $\rightarrow$  This is the tangent space to the circle with vectors of bounded length  $\rightarrow$  cylinder.



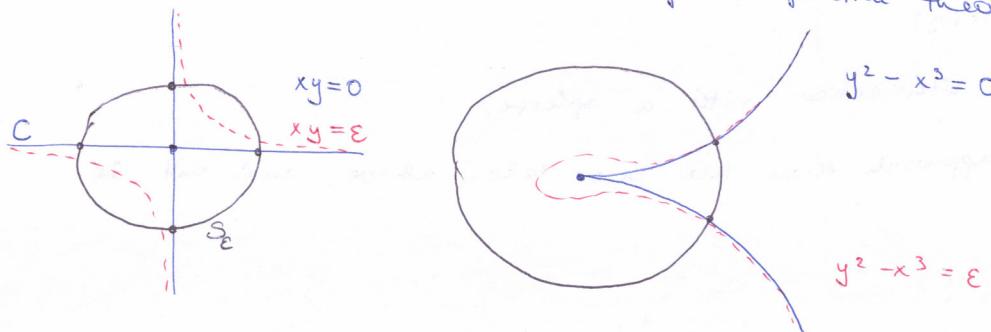
This sufficiently amends the explanation on p. 9. □

Singularities

Def:  $C = \{[x:y:z] \in \mathbb{P}^2 \mathbb{C} \mid f(x,y,z) = 0\}$ .  $p \in C$  is a **singular point** if

$$\frac{\partial f}{\partial x}(p) = \frac{\partial f}{\partial y}(p) = \frac{\partial f}{\partial z}(p) = 0.$$

$p \in C$  is a **smooth point** if at least one of these partial derivatives doesn't vanish (and then the implicit function theorem can be applied).



looking at level sets: for small  $\epsilon$ ,  $f = \epsilon$  is close to  $f = 0$  except at the singular pt, where it is relatively far away.

Intersect with a sphere and study the intersections.

$$S_\epsilon := \{(x,y) \in \mathbb{C}^2 \mid |x|^2 + |y|^2 = \epsilon^2\} \quad \text{3-dimensional sphere}$$

$$C = \{(x,y) \in \mathbb{C}^2 \mid f(x,y) = 0\} \ni (0,0) \quad (\text{translate so that this holds})$$

Study  $S_\epsilon \cap C$ .

$$\text{Ex: } f(x,y) = x^2 - y^3.$$

Write  $x = r e^{i\alpha}$ ,  $y = s e^{i\beta}$  polar coordinates,  $r, s > 0$ ,  $\alpha, \beta \in \mathbb{R}$

$$\begin{aligned} (x,y) \in S_\epsilon \cap C &\iff s^2 e^{2i\beta} = r^3 e^{3i\alpha} \quad \& \quad s^2 + r^2 = \epsilon^2 \\ &\iff s^2 = r^3 \quad \& \quad 2\beta - 3\alpha \in 2\pi\mathbb{Z} \quad \& \quad \underbrace{r^3 + s^2}_{r^3 + r^2} = \epsilon^2 \end{aligned}$$

strictly mon increasing

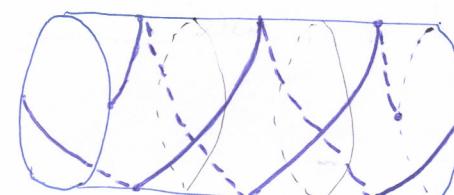
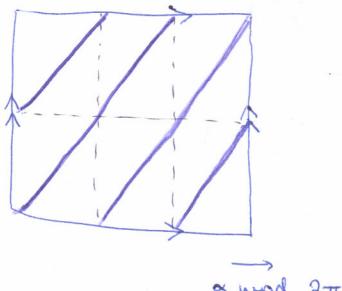
$\Rightarrow$  has unique solution  $r$

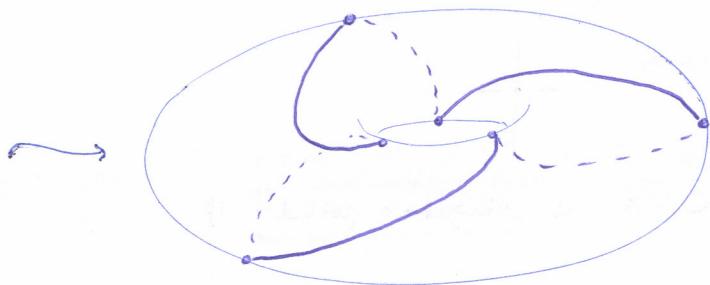
This is a torus.

$$\Rightarrow \text{get unique } s = r^{3/2} \Rightarrow (x,y) \in \underbrace{r S^1 \times s S^1}_{\subset S_\epsilon^3 \subset \mathbb{C}^2}$$

$$\beta = \frac{3}{2}\alpha + n\pi$$

$$n \in \mathbb{Z}$$





We get a trefoil knot.

(For the time being, we don't care about how these are oriented, positivity, and negativity.)

Now we take a different approach: parametrize the curve:

$$C \rightarrow C \subset \mathbb{C}^2$$

$$t \mapsto (t^2, t^3) = (x, y)$$

For  $|t|$  fixed, we get the intersection with a sphere.

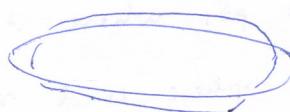
This is a more general approach than the one taken above, and will be generalized later.

Ex.  $f = y^2$

$$(x, y) \in S_\varepsilon \cap C \Leftrightarrow |x| = \varepsilon, y = 0$$



perturb,  
every point splits  
into 2 points



there is some twisting here

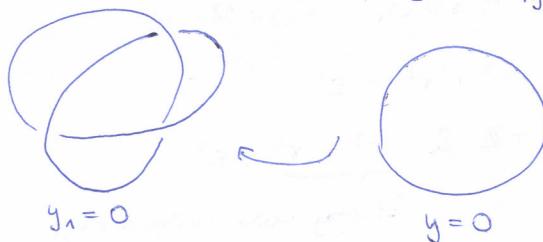
$$\text{Ex. } f = (\underbrace{y^2 - x^7}_y)^2 - x^7 = y_1^2 - x^7$$

We start with this because this term has minimal degree.

$\forall x \neq 0$  there are exactly 2  $y$ 's close to  $y_1 = 0$ :

The  $x^7$  in comparison is much smaller.

$$y_1 = x^{7/2}, \quad x \text{ small, } |x| \approx \varepsilon \Rightarrow |y_1| \approx \varepsilon^{7/2} \ll \varepsilon$$



This motivates the following procedure.

Def. A polynomial  $f \in \mathbb{C}[x, y]$  is **weighted homogeneous** if  $\exists n, a, b \in \mathbb{N}$  s.t.

$$f(t^a x, t^b y) = t^n f(x, y) \quad \forall t \in \mathbb{C} \quad \forall (x, y) \in \mathbb{C}^2$$

$w(f) := n$  weight of  $f$ ,  $w(x) := a$ ,  $w(y) := b$  weights of  $x$  resp.  $y$

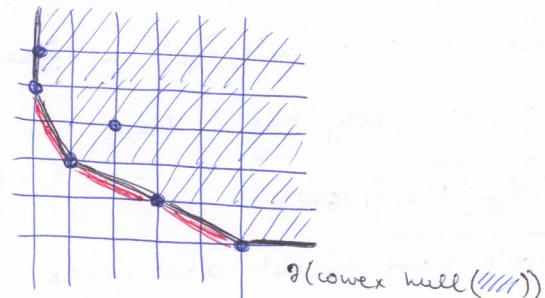
Ex.  $f(x,y) = y^2 - x^3 \rightarrow w(x)=2, w(y)=3, w(f)=6$  weighted homogeneous

$x^i y^j$  is also weighted homogeneous,  $w(x)=j, w(y)=i, w(x^i y^j)=2ij$

Def.  $f = \sum_{i,j} a_{ij} x^i y^j \in \mathbb{C}[x,y]$  has support  $\text{supp } f = \{(i,j) \in \mathbb{N}^2 \mid a_{ij} \neq 0\}$

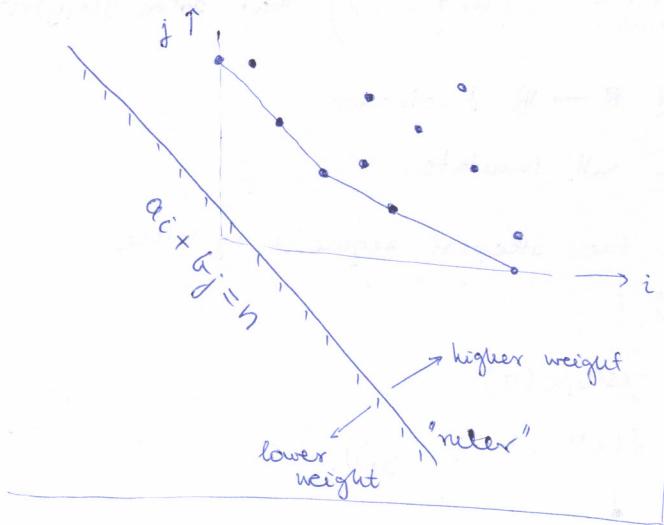
Note that this is of course a notion totally different from that of a support of a fraction.

Ex.  $f = 7y^5 - 5y^4 + xy^2 - 3x^2y^3 + x^3y + x^5$



$\partial(\text{convex hull} \left( \bigcup_{(i,j) \in \text{supp } f} p + \mathbb{R}_{>0}^2 \right))$  is a polygon with 2 infinite rays

The **Newton polygon** of  $f$  consists of the finite segments.



**Theorem (Puiseux).** Let  $C \subset \mathbb{C}^2$  be an algebraic curve. Assume that  $(0,0) \in C, \{x=0\} \not\subset C$ . Then near  $(0,0)$   $C$  is the union of finitely many branches  $\gamma_1, \dots, \gamma_n \subset \mathbb{C}^2$  where each branch has an injective parametrisation of the form  $t \mapsto (t^m, g(t))$  where  $m \in \mathbb{N}_1$ ,  $g$  a holomorphic function (i.e. a power series in  $t$ ).

Moreover, the branches intersect at  $(0,0)$  only. fractional powers is called a **Puiseux series**.

Rewrite  $g(x^{1/m})$ : a **Newton series** is

$$x^{\frac{q_1}{p_1}} \left( a_1 + x^{\frac{q_2}{p_1 p_2}} \left( a_2 + x^{\frac{q_3}{p_1 p_2 p_3}} \left( \dots \right) \right) \right)$$

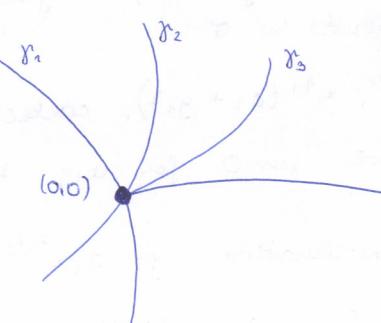
where  $\mathbb{A}(p_i, q_i)$  is a **Newton pair**, i.e.

these are coprime positive integers.

$$b_1 x^{\frac{m_1}{n_1}} + b_2 x^{\frac{m_2}{n_1 n_2}} + b_3 x^{\frac{m_3}{n_1 n_2 n_3}}$$

where  $\mathbb{A}(m_i, n_i)$  is a **Puiseux pair**; i.e. coprime

$p_i := n_i, q_i := m_i, q_i := m_i - m_{i-1} n_i$  gives a bijection between these



Goal for today: "solve"  $f(x,y) = 0$  near  $(0,0)$ ;  $f(0,0) = 0$ .

- If  $\frac{\partial f}{\partial y}(0,0) \neq 0 \rightarrow$  the implicit function theorem does the job.

The implicit function theorem gives a holomorphic output if the input is holomorphic  $\Rightarrow$  we get  $y = \sum_{i \geq 0} a_i x^i$  power series

- If  $\frac{\partial f}{\partial x}(0,0) = \frac{\partial f}{\partial y}(0,0) = 0$ , Puiseux's theorem does the job.

Consequence: we only get fractional power series  $y = \sum_{i \geq 0} a_i x^{i/N}$ , and not only one: there are multiple branches.

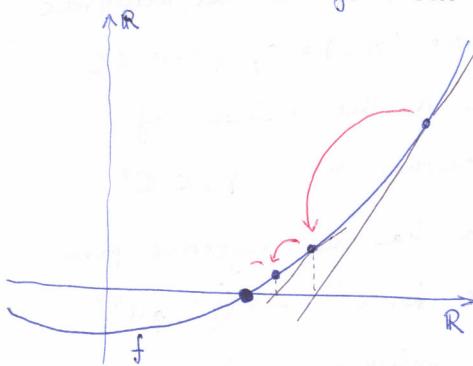
### Newton's Algorithm. (1676)

We won't worry about convergence (Newton for sure didn't); perhaps this will be discussed in an upcoming lecture.

Input:  $f \in \mathbb{C}[x,y] \setminus \{0\}$ ,  $f(0,0) = 0$

Output: all possible Newton series  $y = x^{q_1/p_1} (a_1 + x^{q_2/p_2} (a_2 + \dots))$  that solve  $f(x,y) = 0$ .

Recall Newton's algorithm for finding zeros of  $\mathbb{R} \rightarrow \mathbb{R}$  functions.



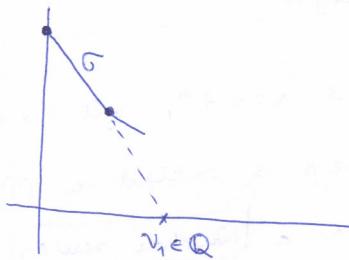
This is what we will imitate.

[STEP 1.] Let  $\sigma$  be the steepest segment of the Newton polygon of  $f$

[STEP 2.]  $\frac{q_1}{p_1} := -1/\text{slope}(\sigma)$

[STEP 3.] Compute  $f(x^{p_1}, x^{q_1} (a_1 + y_1))$ .

$$f(x,y) = \underbrace{\sum_{p_i+q_j=p_1v_1} a_{ij} x^i y^j}_{\text{quasi-homogeneous, corresponds to } \sigma} + \underbrace{\sum_{p_i+q_j>p_1v_1} a_{ij} x^i y^j}_{\text{higher order terms}}$$



In  $f(x^{p_1}, x^{q_1} (a_1 + y_1))$ , collect lowest order terms in  $x$  alone, and solve  $f = 0$  for  $a_{ij}$  these may be several solutions.

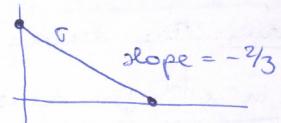
First approximation:  $y = a_1 x^{p_1/q_1} + \dots$

[STEP 4.]  $f_1(x,y_1) := f(x^{p_1}, x^{q_1} (a_1 + y_1)) \cdot x^{-p_1 v_1} \in \mathbb{C}[x,y]$  by construction, this is the highest power of  $x$  that can be factored out

[STEP 5.] Repeat for  $f_1$  instead of  $f$ . Obtain  $q_2/p_2, a_2, f_2, \dots$

Instructive example:  $f(x, y) = y^2 - x^3$

$$\frac{q_1}{p_1} = \frac{3}{2} \quad f(x^2, x^3(a_1 + y_1)) = x^6(a_1 + y_1)^2 - x^6 = x^6(a_2 - 1) + \text{higher order terms}$$



The only case when this doesn't work is when  $y_1 \nmid f(x, y_1)$ .

This means that the Newton polygon is above some horizontal line.

The different choices for  $a_i$  give the branches.

Note that the algorithm may be performed with any slope  $\sigma$  instead of the steepest one; Filip chose close to present this version

only out of cosmetic reasons.

More sophisticated example:  $f(x, y) = (y^2 - x^3)^2 - 4x^5y$

$$\frac{q_1}{p_1} = \frac{3}{2} \quad y = a_1 x^{3/2} + \dots$$

$$\begin{aligned} f(x^2, x^3(a_1 + y_1)) &= (x^6(a_1 + y_1)^2 - x^6)^2 - 4x^{13}(a_1 + y_1) - \\ &= x^{12} \cdot \underbrace{(a_1^2 - 1)^2}_{+ \text{higher order terms}} + \text{higher order terms} \\ &= 0 \quad \rightarrow a_1 = \pm 1, \text{ these give us two branches} \end{aligned}$$

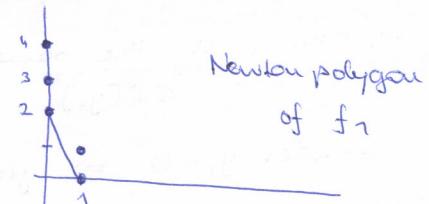
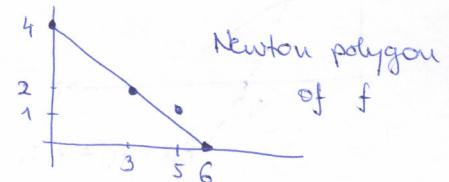
$$\begin{aligned} v_1 = 6, \quad f_1(x, y_1) &= x^{-12} f(x^2, x^3(\pm 1 + y_1)) \\ &= (\pm 2y_1 + y_1^2) - 4x(\pm 1 + y_1) \end{aligned}$$

$$\frac{q_2}{p_2} = \frac{1}{2}$$

$$f_1(x^2, x^1(a_2 + y_1)) = (\pm 2x(a_2 + y_2) + x^2(a_2 + y_2)^2)^2 - 4x^2 - 4x^3(a_2 + y_2)$$

$$\Rightarrow a_2 = \begin{cases} \pm 1 & \text{if } a_1 = 1 \\ \pm i & \text{if } a_1 = -1 \end{cases} \quad = 4x^2 a_2^2 + 4x^2 + \text{higher order terms}$$

$$y = x^{3/2} \left( \pm 1 + x^{1/2 \cdot 2} \left( \frac{\pm 1}{\pm i} + \dots \right) \right)$$



Claim. The denominators are bounded, i.e. Newton's algorithm gives a series in  $x^{1/N}$  for some  $N \in \mathbb{N}$ .

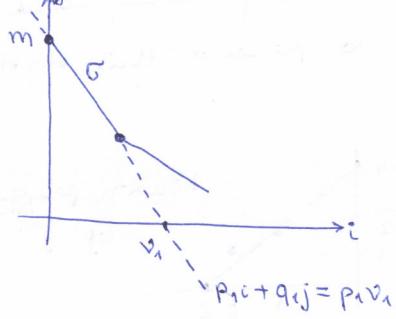
$$\text{PF: } y = x^{q_1/p_1} (a_1 + x^{q_2/p_2} (a_2 + \dots))$$

We will show that  $\exists i_0: p_i = 1 \forall i \geq i_0$ . Then take  $N := p_1 \cdots p_{i_0}$ .

$$\text{Recall that } f_1(x, y_1) = x^{-p_1 v_1} \cdot f\left(x^{p_1}, x^{q_1} (a_1 + y_1)\right)$$

$$\sum_{i p_1 + j q_1 = p_1 v_1} a_{ij} x^{ip_1} y_1^{jq_1} (a_1 + y_1)^j + \sum_{i p_1 + j q_1 > p_1 v_1} \dots$$

Consider the steepest slope  $\sigma$  of the Newton polygon.



$$m = \frac{p_1}{q_1} v_1 \quad f(0, y) = y^m \cdot \text{const} + \text{l.o.t.}$$

$$\begin{aligned} f_1(0, y) &= \sum_{i p_1 + j q_1 = p_1 v_1} a_{ij} (a_1 + y_1)^j + 0 \\ &= \text{const} \cdot y_1^{\frac{p_1 v_1}{q_1}} + \text{lower order terms in } y_1 \end{aligned}$$

So if  $m_1$  denotes the largest exponent of  $y_1$  in  $f_1(0, y)$  then  $m_1 \leq \frac{p_1 v_1}{q_1} = m$

By induction we obtain  $m \geq m_1 \geq m_2 \geq \dots$

When do we have equality? It suffices to study  $m_1 = m$ .

$$f_1(0, y_1) = g(a_1 + y_1) \text{ for some polynomial } g(t) = \sum_{i p_1 + j q_1 = p_1 v_1} a_{ij} t^j \in \mathbb{C}[t]$$

$$\text{Consider } y_1 = 0 \Rightarrow g(a_1) = 0 \quad (a_1 \neq 0 \text{ by def})$$

$$\deg g = m = \frac{p_1 v_1}{q_1}$$

$$\begin{aligned} m_1 &= (\text{order of the zero } y_1 = 0 \text{ of } g(a_1 + y_1) \in \mathbb{C}[y_1]) \\ &= (\text{order of the zero } a_1 \text{ of } g(t) \in \mathbb{C}[t]) \end{aligned} \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{equal by assumption}$$

$$\Rightarrow (\text{const} \neq 0) \cdot (t - a_1)^m = g(t) = \sum_{i p_1 + j q_1 = p_1 v_1} a_{ij} t^j$$

coeff of  $t^{m-1}$  is nonzero  $\Rightarrow a_{i, m-1} \neq 0$  for some  $i$

$$\text{s.t. } p_1 i + q_1(m-1) = p_1 v_1$$

$$\Rightarrow p_1 i + p_1 v_1 - q_1 = p_1 v_1 \Rightarrow i = \frac{q_1}{p_1}$$

But  $i \in \mathbb{N}$  and  $(p_1, q_1) = 1 \Rightarrow p_1 = 1$

Thus  $p_1 \dots p_{i+1}$  either does not increase (when  $p_{i+1} = 1$ )  
or  $m_{i+1} < m_i$ .

## Knots, links & surfaces

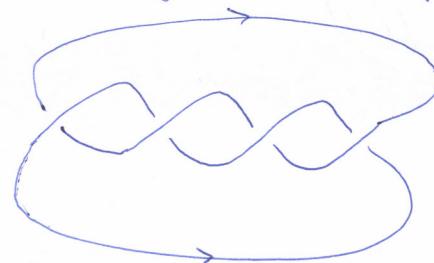
14.05.2019

This will be a non-comprehensive, non-precise introduction, just to give us some vocabulary to work with.

Def. **Knot**: oriented, smoothly embedded  $S^1 \hookrightarrow S^3$ .

$$S^3 = \mathbb{R}^3 \cup \{\infty\} = \{(x,y) \in \mathbb{C}^2 \mid |x|^2 + |y|^2 = 1\} = \{p \in \mathbb{R}^4 \mid \|p\| = 1\}$$

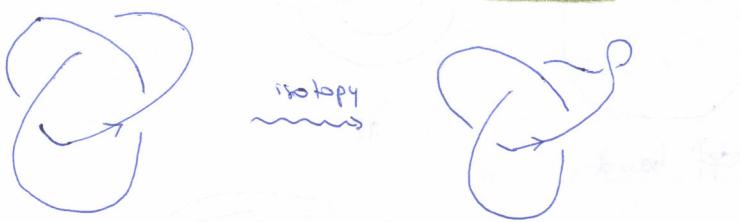
Def. **Link**: disjoint union of fin many knots:  $S^1 \amalg \dots \amalg S^1 \hookrightarrow S^3$



Def. **Isotopy**:  $\varphi: [0,1] \times S^3 \rightarrow S^3$  continuous,

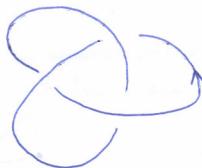
$\varphi_t(p) := \varphi(t, p)$ ;  $\varphi_t: S^3 \rightarrow S^3$  is a homeomorphism  $\forall t \in [0,1]$ ,  $\varphi_0 = \text{id}_{S^3}$

Def.  $K_0, K_1 \hookrightarrow S^3$  knots are **isotopic** if  $\exists \varphi$  isotopy s.t.  $\varphi_1(K_0) = K_1$ .



$$\begin{array}{ccc} K \hookrightarrow S^3 \setminus \{\text{point}\} & \cong & \mathbb{R}^3 \\ & \searrow \text{projection} & \\ & \Rightarrow \mathbb{R}^2 & \end{array}$$

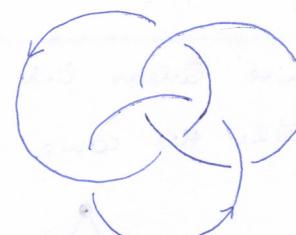
We can do this so that there are no triple crossings, every crossing looks like 

Ex.

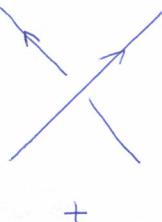
Trefoil knot



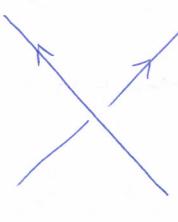
Figure 8 knot



Borromean rings

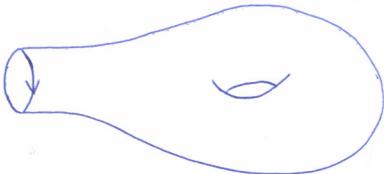


Positive crossing

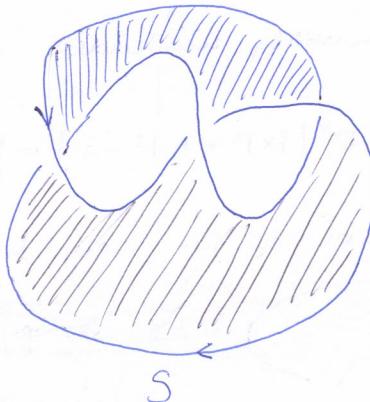


Negative crossing

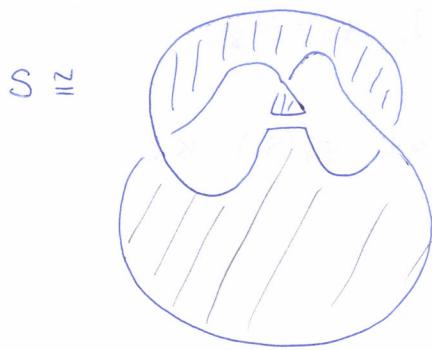
Def. **Seifert surface**: compact, connected, oriented surface with boundary  $S \hookrightarrow S^3$ .



Seifert surface that has the unknot as its boundary



These surfaces are homeomorphic:



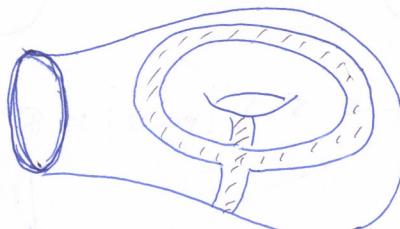
$\approx$



$\approx$



Hopf band



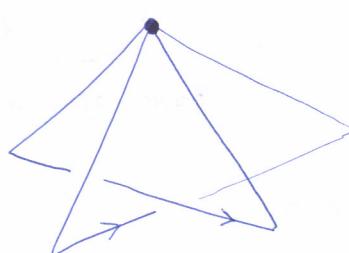
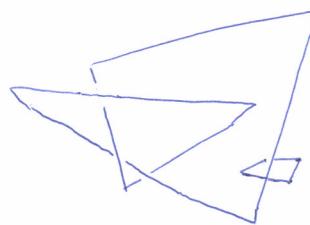
$\approx$

But they are not isotopic since their boundaries are not isotopic (let's just believe that for now.)

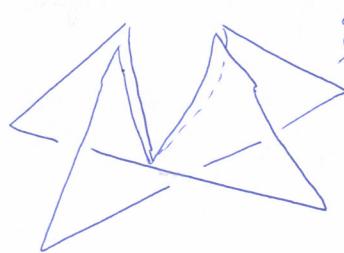
Thm. (Frankl-Pontryagin, Seifert)  
1929 1934  $VL \hookrightarrow S^3$  link  $\exists S \hookrightarrow S^3$  Seifert surface:  $\partial S = L$ .

Pf: Consider a piecewise linear link "near" a plane

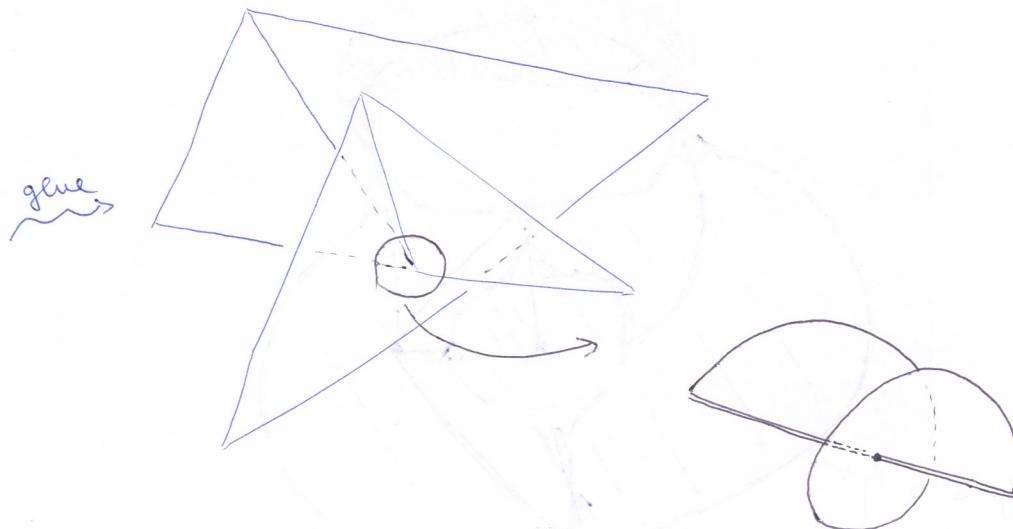
Take the cone  $\rightarrow$  obtain a cell complex.



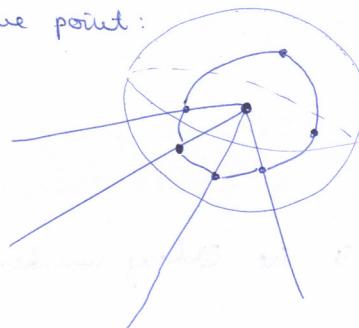
cut  
off the tip



glue

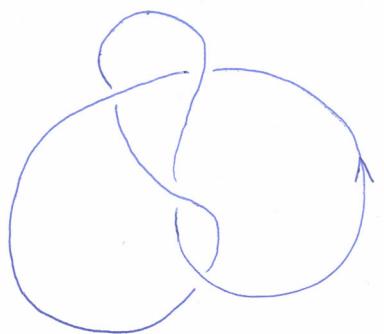


Near the cone point:

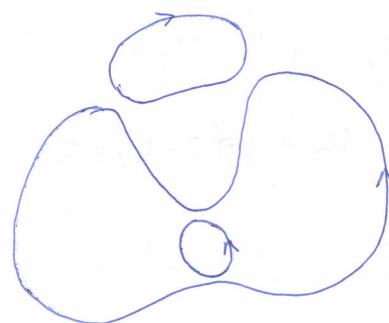


remove the ball and glue in disjoint discs.

Seifert's PROOF for the Thm. through an example:

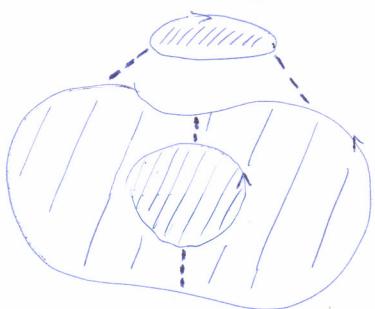


Always replace  $\nearrow$  and  $\nwarrow$  by  $\nearrow$

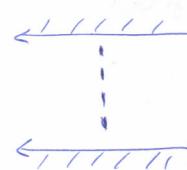


Thus we obtain  
a disjoint union  
of oriented circles.

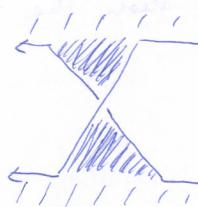
We place each circle on a separate plane, these planes being above one another. Use Jordan's Thm on each of these planes.

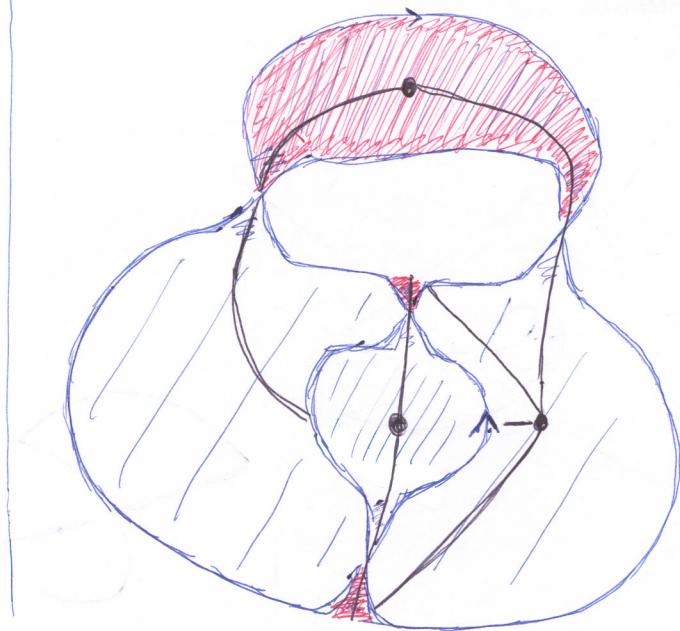
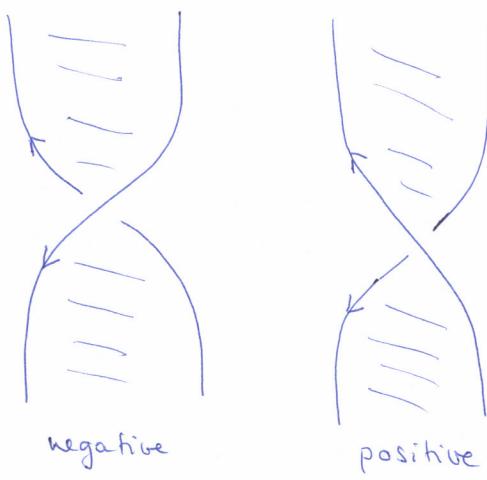


For each former crossing, glue in a half-twisted band:



now



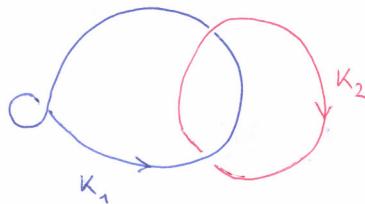


## Linking number

Let  $K_1, K_2 \subset S^3$  be disjoint knot

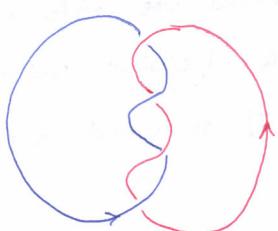
Def:  $\text{lk}(K_1, K_2) := \frac{1}{2} \left( \# \left( \begin{array}{c} \nearrow \searrow \\ K_i & K_j \\ i \neq j \end{array} \right) - \# \left( \begin{array}{c} \nearrow \nearrow \\ K_i & K_j \\ i \neq j \end{array} \right) \right)$  is the linking number of  $K_1, K_2$ .

Ex.



$$\text{lk}(K_1, K_2) = \frac{1}{2} (2 - 0) = 1.$$

Ex.



$$\text{lk} = \frac{1}{2} (0 - 4) = 2$$

Rule.

$$\text{lk}(K_1, K_2) = \# (S \cap K_2) \quad \text{where } K_1 = \partial S.$$

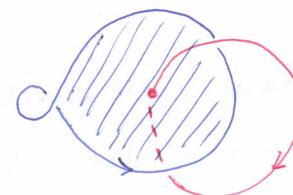
Rule.

$$\text{lk}(K_1, K_2) = \pm [K_2] \in H_1(S^1 \setminus K_1; \mathbb{Z}) \cong \mathbb{Z}$$

$$= \pm [K_1] \in H_1(S^2 \setminus K_2; \mathbb{Z}) \cong \mathbb{Z}$$

choose a generator

Once and for all, then the sign is fixed

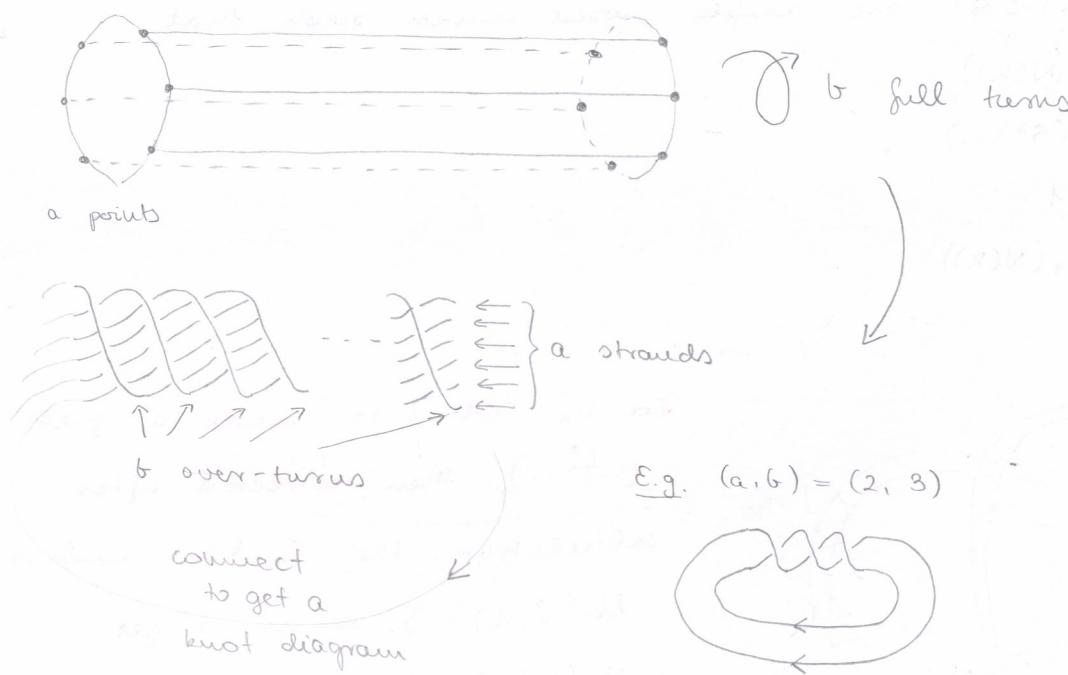


Recall that a torus is a surface homeomorphic to  $S^1 \times S^1$ .

Def. A **torus knot**  $K$  is a knot which can be drawn (i.e. embedded) on the surface of a standardly embedded torus. (into  $S^3$ )

Def. For  $(a, b) \in \mathbb{Z}^2$ ,  $(a, b) = 1$  let  $T(a, b)$  be the torus knot representing the class  $(a, b) \in \mathbb{Z} \oplus \mathbb{Z} \cong H_1(S^1 \times S^1; \mathbb{Z})$ .

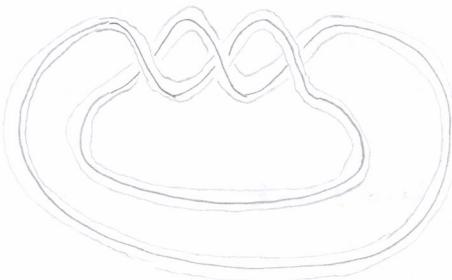
In other words,  $T(a, b)$  wraps around a torus  $a$  times in one direction,  $b$  times in the other.



Recall that homology also sees orientation.

Rule.  $\# \pi_0(\{y^n - x^m = 0\} \cap S^3) = \gcd(n, m)$ . This is why we assume  $(a, b) = 1$ .

Rule.  $T(a, b) = T(b, a)$ , and  $T(\pm 1, b) = unknot$



Take tubular fibres:

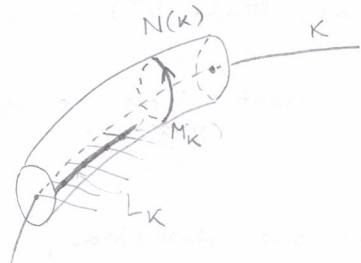
$S^1 \hookrightarrow S^3$  gives rise to  $S^1 \times D^2 \hookrightarrow S^3$

$\text{fundamental group of } S^1 \times D^2 \cong \langle a, b \mid ab^{-1}a^{-1}b \rangle$

## Meridian and longitude

Let  $K \hookrightarrow S^3$  be a knot. Take a tubular nbhd of  $K$  in  $S^3$ :

formally  $K: S^1 \hookrightarrow S^3$  induces  $N(K): S^1 \times D^2 \hookrightarrow S^3$  s.t.  $N(K)(S^1 \times \{0\}) = K(S^1)$



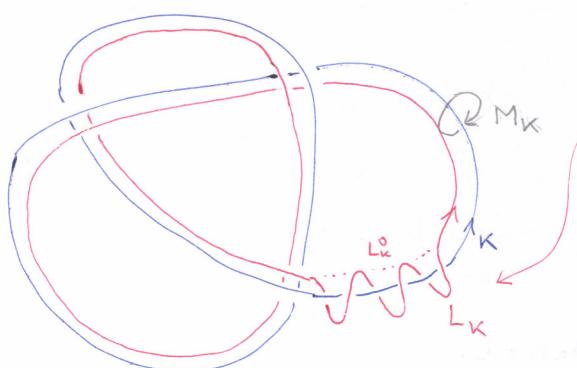
Rule.  $\partial N(K) \cong S^1 \times S^1$ .

Rule. We identify the embeddings  $K, N(K)$  with their images.

Def.  $M_K, L_K \subset \partial N(K) \subset S^3$  are simple closed curves such that

- $[M_K] = 1 \in H_1(S^3 \setminus K)$
- $[L_K] = 0 \in H_1(S^3 \setminus K)$
- $lk(M_K, L_K) = 1$
- $[L_K] = [K] \in H_1(N(K))$

Ex. Trefoil knot



For  $L_K$ , we first made a guess ( $-L_K^\circ$ ), then corrected after calculating the linking numbers  $lk(L_K^\circ, K) = 3$ . so as to get  $lk(L_K, K) = 0$ .

## Cable of a knot

$$K \hookrightarrow S^3, \quad \psi: S^1 \times S^1 \xrightarrow{\cong} \partial N(K)$$

$$S^1 \times \{1\} \longrightarrow L_K$$

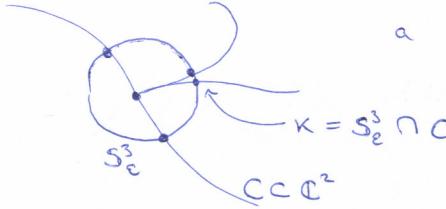
$$\{1\} \times S^1 \longrightarrow M_K$$

Def. For  $(a, b) \in \mathbb{Z}^2$ , the  $(a, b)$ -cable of  $K$  is the knot/link

$$K_{(a,b)} := \psi(T(a,b)) \subset \partial N(K) \subset S^3$$

Explicitly:  $K_{(a,b)} = aL_K + bM_K \in H_1(\partial N(K))$ .

Recall: we wish to describe the intersection of  $S^3_\varepsilon$  with a curve  $C$  around a singularity.



Theorem. Let  $y = x^{q_1/p_1} (a_1 + x^{q_2/p_1 p_2} (a_2 + \dots + a_s x^{q_s/p_1 \dots p_s}) \dots)$  be the Newton series parametrising a branch of an algebraic curve  $C$  near  $(0,0)$ . Then for sufficiently small  $\varepsilon > 0$  the knot  $K = S^3_\varepsilon \cap C$  is the  $(p_s, q_s)$ -cable of the  $(p_{s-1}, q_{s-1})$ -cable of ... of the  $(p_1, q_1)$ -cable of the unknot  $\text{O}$  where  $\alpha_1 := q_1$ ,  $\alpha_{i+1} := q_{i+1} + p_{i+1} p_i \alpha_i \quad \forall i \geq 1$ .

If  $(p_i, \alpha_i) = 1$  s.t.  $p_i > 0$ ,  $\alpha_i > 0$ ,  $\alpha_{i+1} > p_{i+1} p_i \alpha_i$  then we can reconstruct  $(p_i, q_i)$  s.t.  $p_i, q_i > 0$ ,  $(p_i, q_i) = 1$ .

⇒ The knot of the singularity is  $\left( \left( \text{O}_{(p_1, \alpha_1)} \right)_{(p_2, \alpha_2)} \right)_{(p_3, \alpha_3)} \dots \right)_{(p_s, \alpha_s)}$

Pf OF THM: Replace  $S^3_\varepsilon$  by  $R_\varepsilon := \{(x, y) \in \mathbb{C}^2 \mid |x| = \varepsilon, |y| \leq \varepsilon\} \cong S^1 \times \mathbb{D}^2$

(We will explain why this is possible later.)

Consider the  $i$ th approximation of  $y$ :

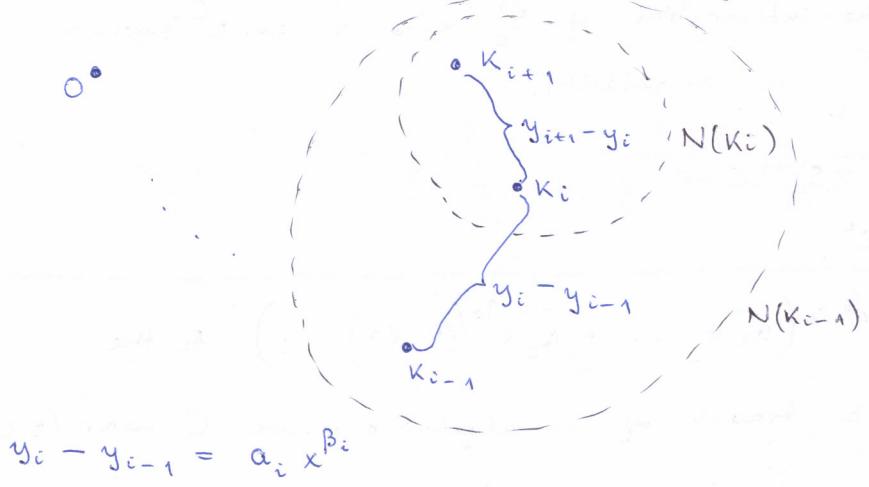
$$y_i = x^{q_1/p_1} (a_1 + x^{q_2/p_1 p_2} (a_2 + \dots + x^{q_i/p_1 \dots p_i}) \dots)$$

This defines a knot  $K_i \subset R_\varepsilon$ .

Induction on  $i$ :

- $y_1 = a_1 x^{q_1/p_1}$  describes the torus knot  $K_1 = T(p_1, q_1)$ ; by an exercise, this is the  $(p_1, q_1)$ -cable of  $\text{O}$ .
- $i$  to  $i+1$ : study  $K_j$  for  $j = i-1, i, i+1$ .
  - × fixed,  $|x| = \varepsilon$ .

The plane of the paper sheet is the  $y$ -disk:



$$y_i - y_{i-1} = a_i x^{\beta_i}$$

$$\beta_i = \frac{q_i}{p_1 \cdots p_s} + \frac{q_{i-1}}{p_1 \cdots p_{i-1}} + \cdots + \frac{q_1}{p_1}$$

$$|y_{i+1} - y_i| = a_{i+1} x^{\beta_{i+1}} \times \frac{q_{i+1}}{p_1 \cdots p_{i+1}} = \underbrace{\left| \frac{a_{i+1}}{a_i} \right|}_{\ll |y_i - y_{i-1}|} |y_i - y_{i-1}| \cdot \underbrace{x^{\frac{q_{i+1}}{p_1 \cdots p_{i+1}}}}_{\text{power of } \epsilon}$$

$\Rightarrow$  we can choose tubular neighborhoods  $N(K_{i-1}), N(K_i)$  such that

$$K_i \subset N(K_i) \subset N(K_{i-1}) \setminus K_{i-1}$$

Parametrise the knot:  $x = \epsilon t^{p_1 \cdots p_i}$ , for  $t \in S^1$

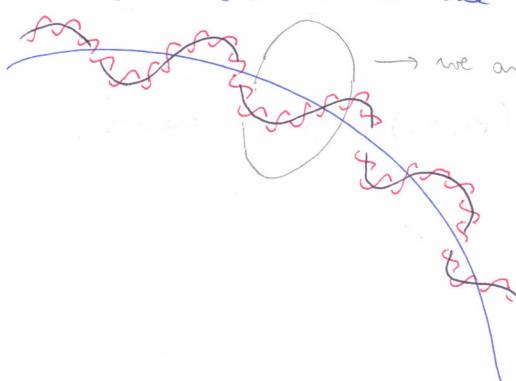
For  $x$  fixed, we get  $p_1 \cdots p_i$  values for  $t$ .

$$y_{i+1} - y_i = \epsilon (y_i - y_{i-1}) \cdot t^{\frac{q_{i+1}}{p_{i+1}}}$$

When  $t$  makes a full turn,  $y_{i-1}, \dots, y_i$  all come back to the same position.

$(y_i - y_{i-1})$  makes some number of full turns;

$(y_{i+1} - y_i)$  makes the same turns as  $(y_i - y_{i-1})$ , plus  $\frac{q_{i+1}}{p_{i+1}}$  of a turn.



$\rightarrow$  we are currently looking at a cross-section of a circular motion around a circular motion around a circular motion ...  
(Think of the Sun, the Earth and the Moon.)

For  $x$  fixed, we have  $p_1 \cdots p_i$  values for  $y_i$ .

For each of these  $y_i$ -values, there are  $p_{i+1}$  values for  $y_{i+1}$ .

Computing the cabling coeffs:

Let  $M_j, L_j$  be the meridian and (preferred, i.e. determined by Seifert surface) longitude of  $K_j$ .

Induction hyp.:  $K_i = p_i L_{i-1} + \alpha_i M_{i-1} \in H_1(N(K_{i-1}) \setminus K_i)$

Let  $L$  be the knot obtained by pushing  $K_i$  away from  $K_{i-1}$  by some small distance  $\delta > 0$ .

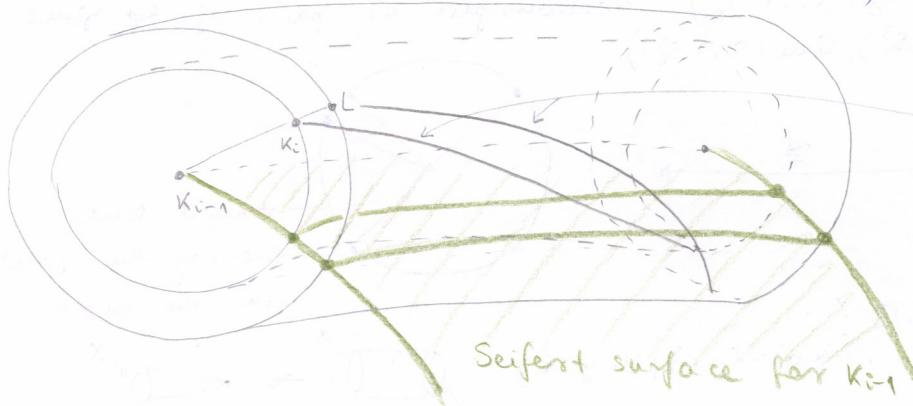
$$y_L = y_i + \delta \times \beta_i$$

$$y_i = y_{i-1} + \alpha_i \times \beta_{i-1}$$

$$K_{i+1} = p_{i+1} L + q_{i+1} M_i$$

$$\text{Wts: } L = L_i + p_i \alpha_i M_i$$

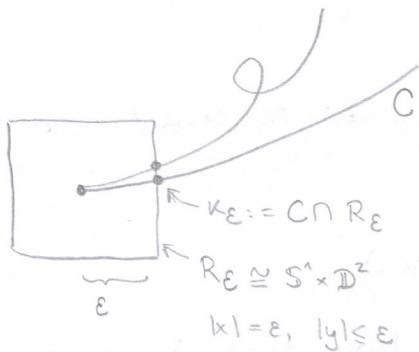
From this:  $K_{i+1} = p_{i+1} L_i + (q_{i+1} + p_{i+1} p_i \alpha_i) M_i$ , which we wanted to prove.



L and  $K_i$   
intersect the  
Seifert surface  
in the same  
way

$$\left. \begin{array}{l} \alpha_i = lk(K_i, K_{i-1}) = lk(L, K_{i-1}) \\ L = K_i \in H_1(N(K_i)) \end{array} \right\} \Rightarrow L = 1L_i + \alpha_i M_{i-1} \in H_1(N(K_{i-1}) \setminus K_i) \\ M_{i-1} = p_i M_i$$

Recall what we did last week:



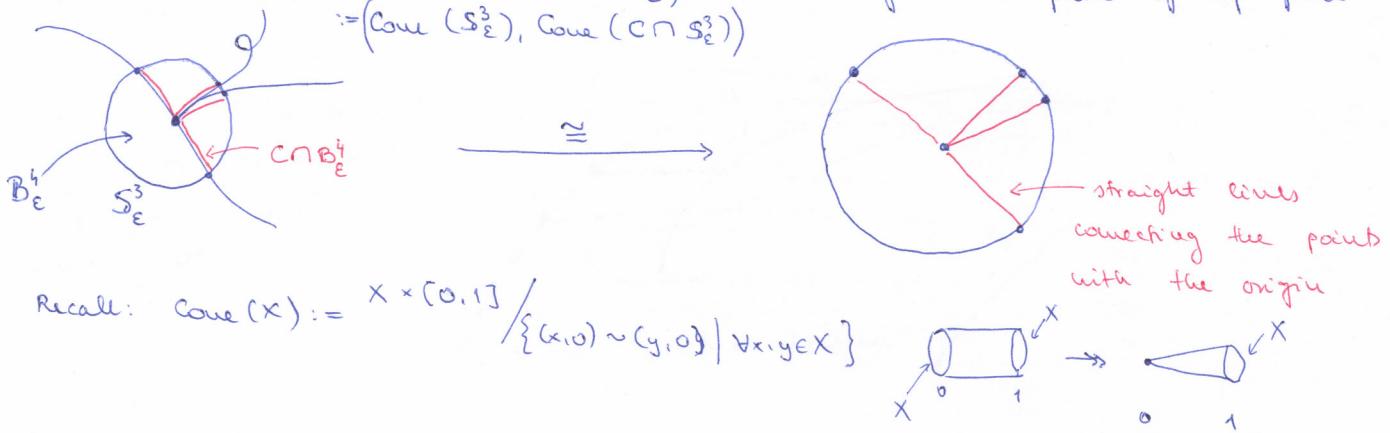
- $K_\varepsilon$  is a link which is independent of  $\varepsilon$  for  $\varepsilon$  sufficiently small.
- $K_\varepsilon$  is an iterated cable of the unknot  $\text{O}$ , with cabling coeffs only depending on the Newton pairs of the branches.

(Tbh we only looked at 1 branch, we didn't investigate what happens when multiple branches are present.)

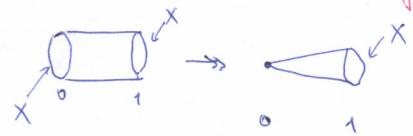
Summary of all this:

Theorem (Conical structure)  $C \subset \mathbb{C}^2$  algebraic curve,  $(0,0) \in C$ . Then  $\exists \varepsilon_0 > 0$

- s.t.  $\forall \varepsilon \leq \varepsilon_0$  s.t. 1)  $K_\varepsilon := C \cap S_\varepsilon^3 \subseteq S_\varepsilon^3 \cong S^3$  is a link indep of  $\varepsilon$   
 2)  $(B_\varepsilon^4, C \cap B_\varepsilon^4) \cong \text{Cone}(S_\varepsilon^3, C \cap S_\varepsilon^3)$  homeomorphic as pairs of top spaces  
 $\cong (\text{Cone}(S_\varepsilon^3), \text{Cone}(C \cap S_\varepsilon^3))$



$$\text{Recall: } \text{Cone}(X) := X \times [0,1] / \left\{ (x,0) \sim (y,0) \mid \forall x,y \in X \right\}$$

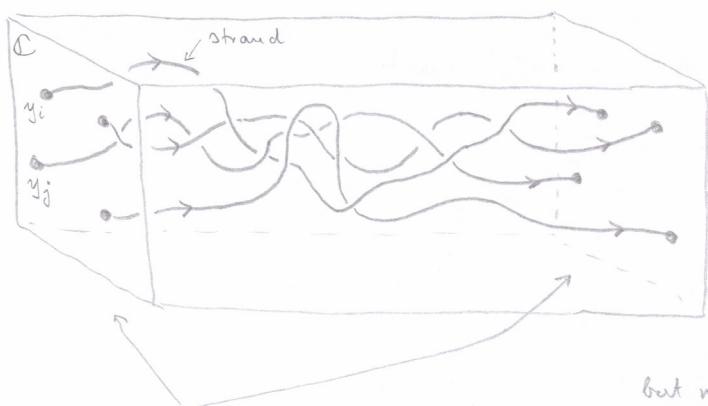


In words: the link of a singularity determines the local embedded topology of the curve  $C$  near the singular point.

Theorem: links of singularities can be represented by positive braid diagrams.

$$P_n := \left\{ p \in \mathbb{C}[y] \mid \deg p = n \text{ and } p \text{ has } n \text{ distinct roots} \right\}$$

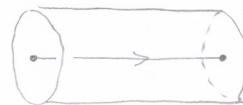
$B_n := \pi_1(P_n)$  braid group on  $n$  strands



but never windide!  
 first and last frame are the same, the points move around, but they need not arrive to the same spot they started from

SKETCH OF PROOF:

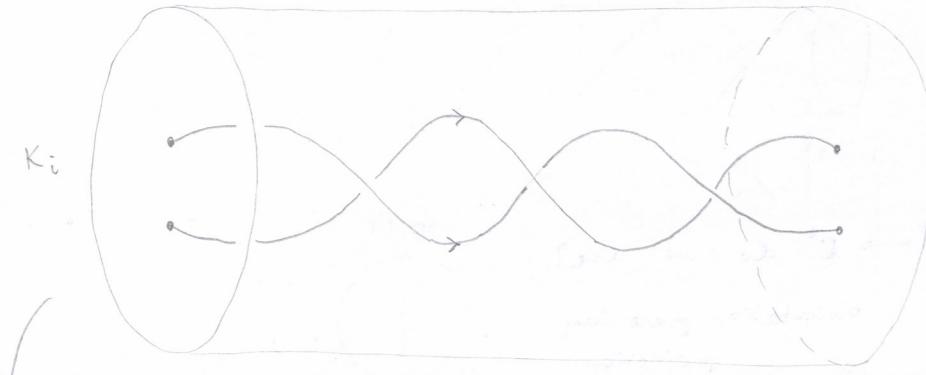
We do induction on the number of cabling operations.



1<sup>st</sup> iteration  $\rightarrow \checkmark$

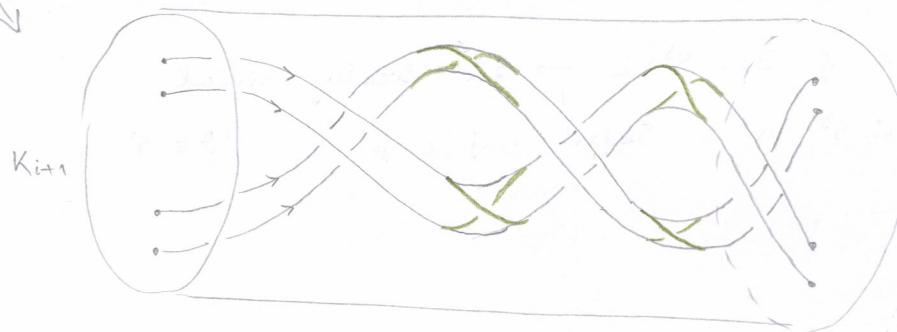
$$y_i - y_{i-1} = a_i \times \beta_i$$

Induction step:



$K_i$ :  $i^{\text{th}}$  iteration of a positive braid

all crossings are +

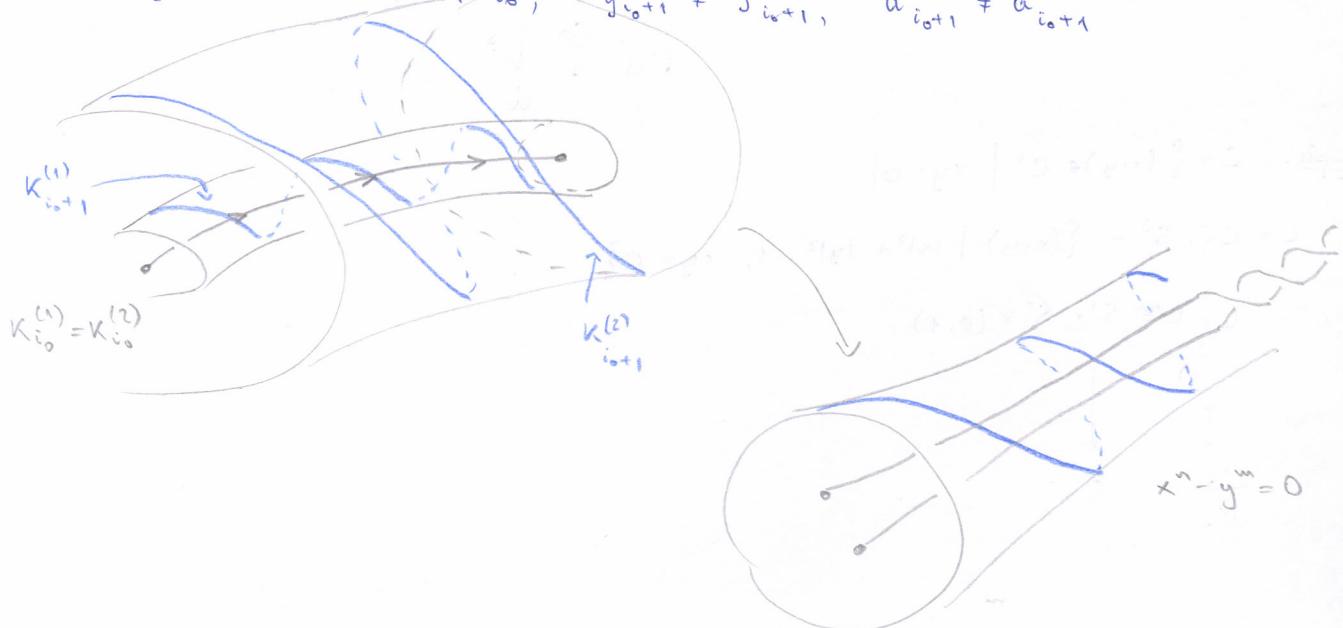


all crossings are still +

For each pos crossing of  $K_i$ , we get  $p_{i+1}^2$  pos crossings for  $K_{i+1}$ , plus additional terms.

$y_i^{(1)}, y_i^{(2)}$ : two Newton series for the two distinct branches

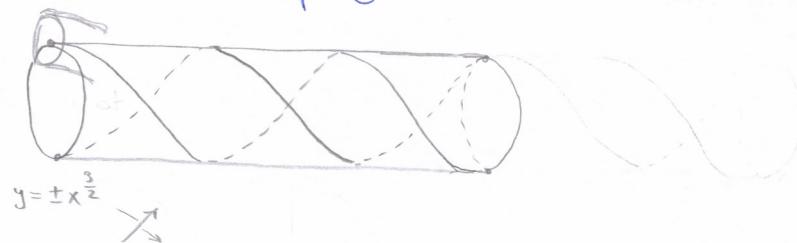
$$y_i^{(1)} = y_i^{(2)} \quad \text{for } i=1, \dots, i_0; \quad y_{i_0+1}^{(1)} \neq y_{i_0+1}^{(2)}, \quad a_{i_0+1}^{(1)} \neq a_{i_0+1}^{(2)}$$



$$\text{Example. } y^{(1)} = x^{\frac{3}{2}} \left( 1 + x^{\frac{1}{4}} \right) = \left( x^{\frac{3}{2}} + x^{\frac{7}{4}} \right) \quad y^{(2)} = x^{\frac{3}{2}} \left( 1 + x^{\frac{1}{6}} \right) = \left( x^{\frac{3}{2}} + x^{\frac{10}{6}} \right)$$

$K^{(1)} := \underbrace{(p_2, \alpha_2)}_{p_2 p_1 x_1 + q_2} - \text{cable of } \underbrace{(p_1, \alpha_1)}_{2} - \text{cable of } \textcircled{0}$

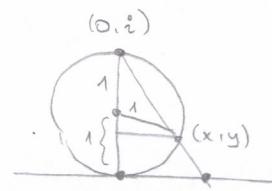
$= (3, 19) - \text{cable of } (2, 3) - \text{cable of } \textcircled{0}$



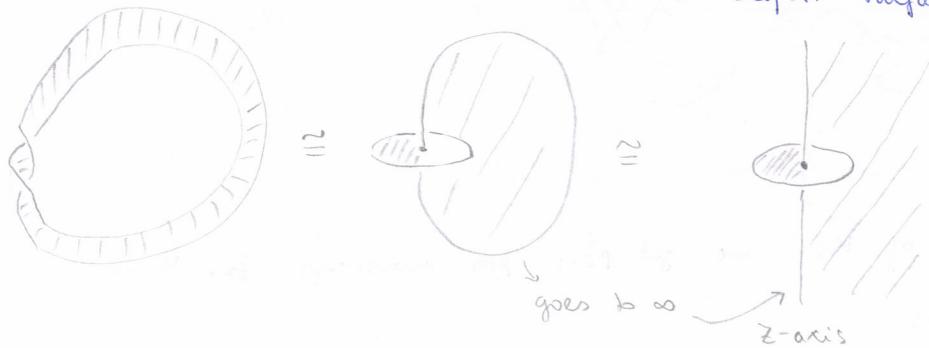
Issue. Which embedding  $S^3 \hookrightarrow \mathbb{C}^2$  do we use?

$$rS^1 \times S^1 \subset S^3 \setminus \{(0, i)\} \longrightarrow \mathbb{R}^3 \text{ orientation-preserving projection}$$

$$(a+ib, c+id) \mapsto \frac{2}{1-d} (a, b, c)$$



Def. A link  $L \subset S^3$  is fibred if  $\exists p: S^3 \setminus L \rightarrow S^1$  locally trivial fibre bundle s.t.  $p^{-1}(s) \cup L \subset S^3$  is a Seifert surface for  $L \quad \forall s \in S^1$ .



Recall.  $p: E \rightarrow B$  is a locally trivial fibre bundle with fibre  $F$  if  $\forall g \in B$   $\exists U \subset B$  open nbh of  $g$  s.t.  $U \times F \cong p^{-1}(U) \subset E$

$$\begin{array}{ccc} & G & \\ p|_U \searrow & \downarrow & \downarrow p \\ & U & \end{array}$$

Example.  $C = \{(x,y) \in \mathbb{C}^2 \mid xy = 0\}$

$$L = C \cap S^3 = \{(x,y) \mid |x|^2 + |y|^2 = 1, xy = 0\}$$

$$S^3 \setminus L \cong S^1 \times S^1 \times (0,1)$$

$$\begin{array}{ccc} (x,y) & & \\ \downarrow & & \downarrow \\ \frac{xy}{|xy|} & & S^1 \end{array}$$

Theo. (Milnor) Links of singularities are fibred:

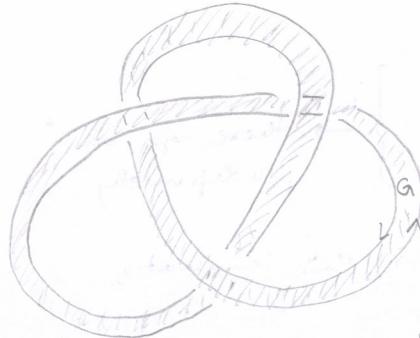
$$L = C \cap S^3_\epsilon, \quad C = \{(x,y) \in \mathbb{C}^2 \mid f(x,y) = 0\}$$

$\frac{f}{|f|} : S^3_\epsilon \setminus L \longrightarrow S^1$  is a locally trivial fibration.

$$q \mapsto \frac{f(q)}{|f(q)|} \quad (f(q) \neq 0 \text{ since } q \notin L)$$

04.06.2019

Exercise 22. The ribbon b/w the two knots is a suitable Seifert surface.

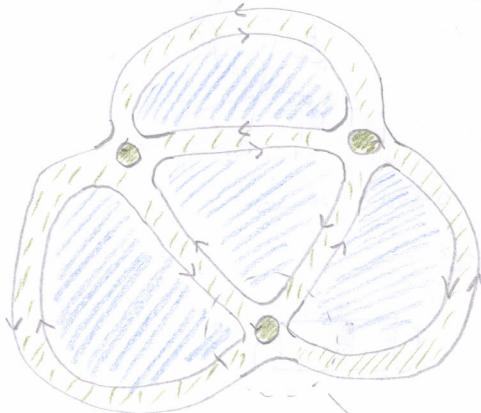


No broken glass can be embedded  $\rightarrow g=0$ .

Or:  $S^2$  has genus 0, and this surface can be embedded into  $S^1$



classification  
of surfaces



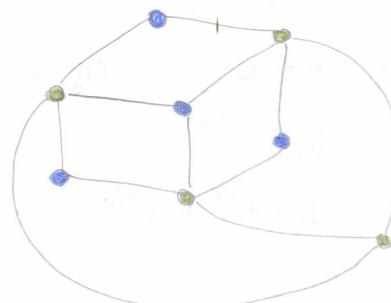
Seifert: replace  $X$  by  $5C$

(Green: counter-clockwise; blue: clockwise)



Then either embed broken glasses to compute the genus,

or use Euler: retract to the graph



$$\left. \begin{aligned} \chi &\stackrel{\text{def}}{=} 8 - 12 = -4 \\ \chi &= 2 - 2g - \frac{b}{2} \end{aligned} \right\} \Rightarrow g = 2$$

Question. Does the Newton algorithm always terminate?

Recall: we stop only if  $y_n \neq f_0$ .

This may never happen. But this need not trouble us: for  $i$  large enough,  $p_i = 1$ , and taking the cable is just an isotopy, hence the knot type does not change.

Note: we have convergence due to WPT.

Answers to 2(a):  $y = x^{3/2} + x^{5/3}$

$$b): y = -2x^2 - 16x^3 - 224x^4 - 3840x^5 - \dots$$

$$y = -\frac{1}{2}x + x^{3/2} + x^2 + \frac{5}{4}x^{5/2} + \dots$$

$$y = -\frac{1}{2}x - x^{3/2} - x^2 + \frac{13}{4}x^{5/2} - \frac{27}{4}x^3$$

→ these go on indefinitely

→ this terminates

a) is from [Ghys], b) from [arXiv: 0807.4674]

Addendum to the Blm on p.25

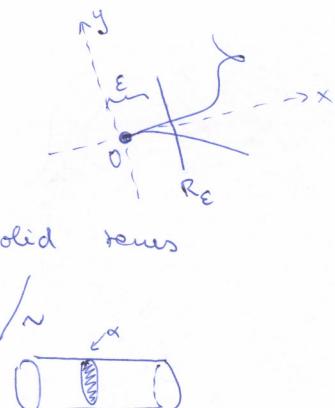
$C \subset \mathbb{C}^2$  branch of an alg curve

$$C \cap R_\varepsilon \cong C \cap S_\varepsilon^3$$

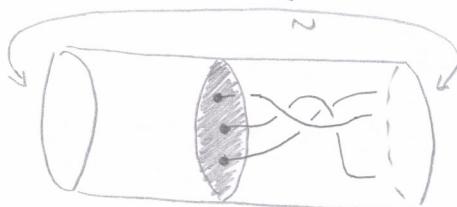
$$R_\varepsilon = \{ |x| = \varepsilon, |y| \leq \varepsilon \} \subset \mathbb{C}^2 \quad \text{solid torus}$$

$$R_\varepsilon \xrightarrow{\cong} S^1 \times D^2 = [0,1] \times D^2 / \sim$$

$$(x,y) \mapsto \left( \frac{x}{\varepsilon}, \frac{y}{\varepsilon} \right)$$



$$x = \varepsilon e^{2\pi i \alpha}, \text{ fixing } x \text{ means fixing } \alpha$$



$$y_1 = a_1 x^{q_1/p_1}$$

$$y_2 = x^{q_1/p_1} \left( a_1 + x^{q_2/p_1 p_2} a_2 \right)$$

We give an example of how this goes.

$$y = x^{3/2} + x^{7/4}$$

→ approximations:

$$y = 0$$

(if  $x$  is small,  $y$  is even smaller)

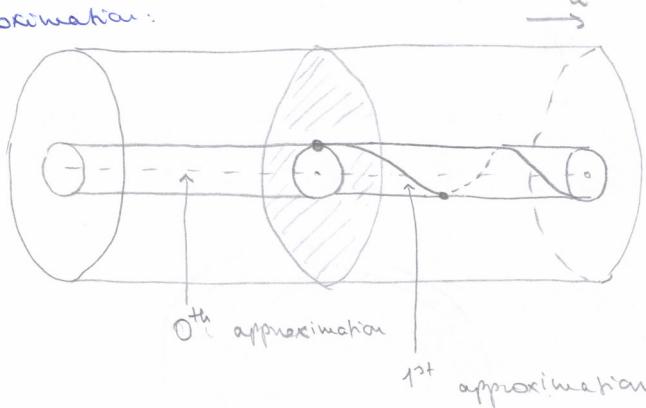


$$x = \varepsilon e^{2\pi i \alpha}$$

$$y_1 = x^{3/2}$$

$$y_2 = x^{3/2} + x^{7/4}$$

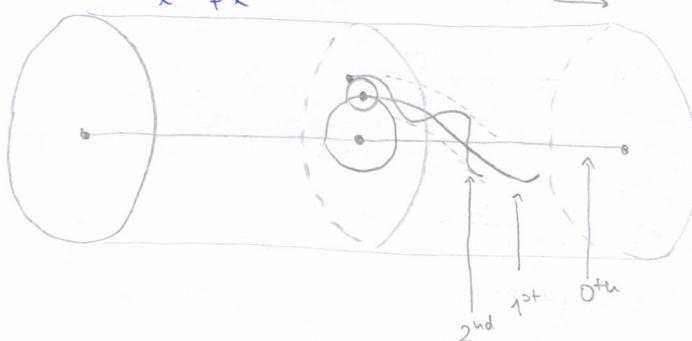
0th approximation

1<sup>st</sup> approximation:

→ we get a torus knot

The error is much compared to the previous step.

$$y_2 = y_1 + \frac{x^{7/4}}{x^{3/2} + x^{1/2,2}}$$

How many times we have to go around to get back to the same point is encoded in the exponents of  $x$ .

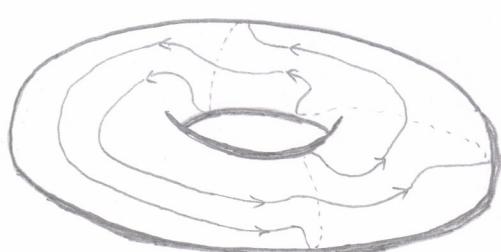
$$y = x^{3/2} \left( 1 + x^{1/2,2} \right)$$

$$\frac{q_1}{p_1} = \frac{3}{2}, \quad \frac{q_2}{p_2} = \frac{1}{2}$$

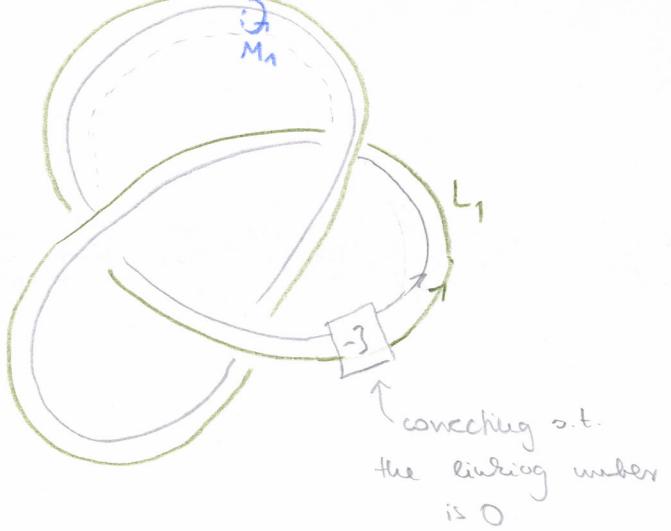
$$\alpha_1 = q_1 = 3 \quad \text{and} \quad \alpha_2 = q_2 + p_2 p_1 x_1 = 13$$

$$K = C \cap R_\varepsilon = \left( \alpha_{(2,3)} \right)_{(2,13)}$$

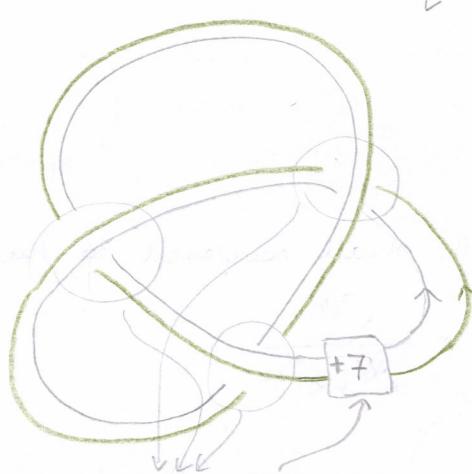
Quark

intersection with the torus:  $L_0$ Seifert surface for  $\alpha$  given by  $\alpha$  $(2,3)$  $\approx$ 

B



Want to draw  $2L_1 + 13M_1$ .



Theorem. Algebraic knots are fibred.

25.06.2019

Recall:  $K = S^3 \cap C$  algebraic knot

$C \subset \mathbb{C}^2$  algebraic curve,  $S^3$  a small sphere

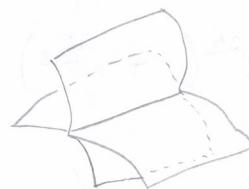
$K$  is fibred if  $\exists p: S^3 \setminus K \rightarrow S^1$  fibre bundle s.t.  $F_t := p^{-1}(t) \cup K$  is a Seifert surface  $\forall t \in S^1$

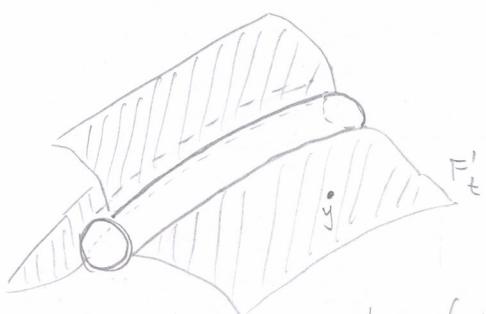
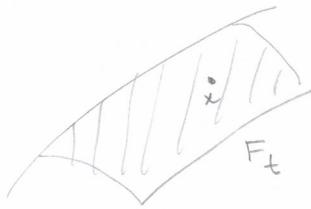
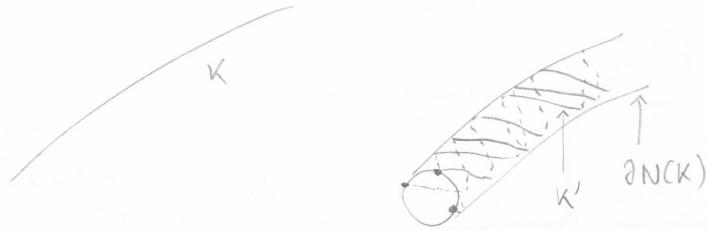
Pf: We know that alg knots are iterated cables of  $\text{O}$ .

We will show the general statement that

if  $K$  is fibred,  $p: S^3 \setminus K \rightarrow S^1$  the fibre bundle s.t.  $\forall t \in (0, 1]: p^{-1}(t) \cup K = F_t$  is Seifert then  $K' := K_{(a/b)}$  is fibred  $\forall a, b \in \mathbb{Z} \setminus \{0\}$ .

That is, the cable of a fibred knot is fibred.





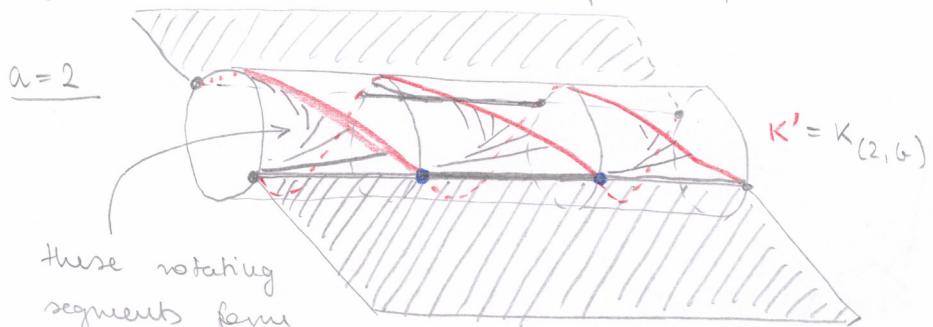
$$p(x) = e^{2\pi i t}$$

For  $y \in S^3 \setminus N(K)$ , define  $p'$  by:

$$p'(y) = (p(y))^a = e^{2\pi i t}, \quad t \in S^1 \setminus K$$

$p'$  will be a map  $S^3 \setminus K' \rightarrow S^1$ ; so far we have defined it only outside  $N(K)$ .

For  $y \in N(K)$ , we won't have an explicit equation.



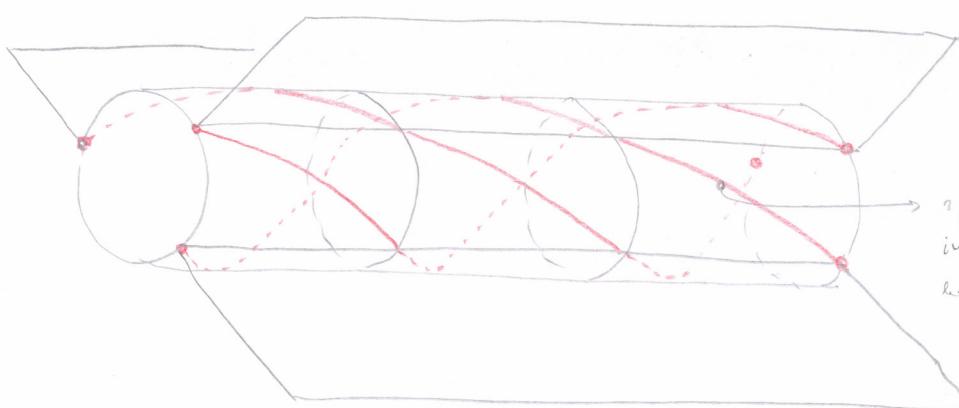
these rotating  
segments form

the fibre bundle  $\rightarrow$  disjoint surfaces except for the gluing pts

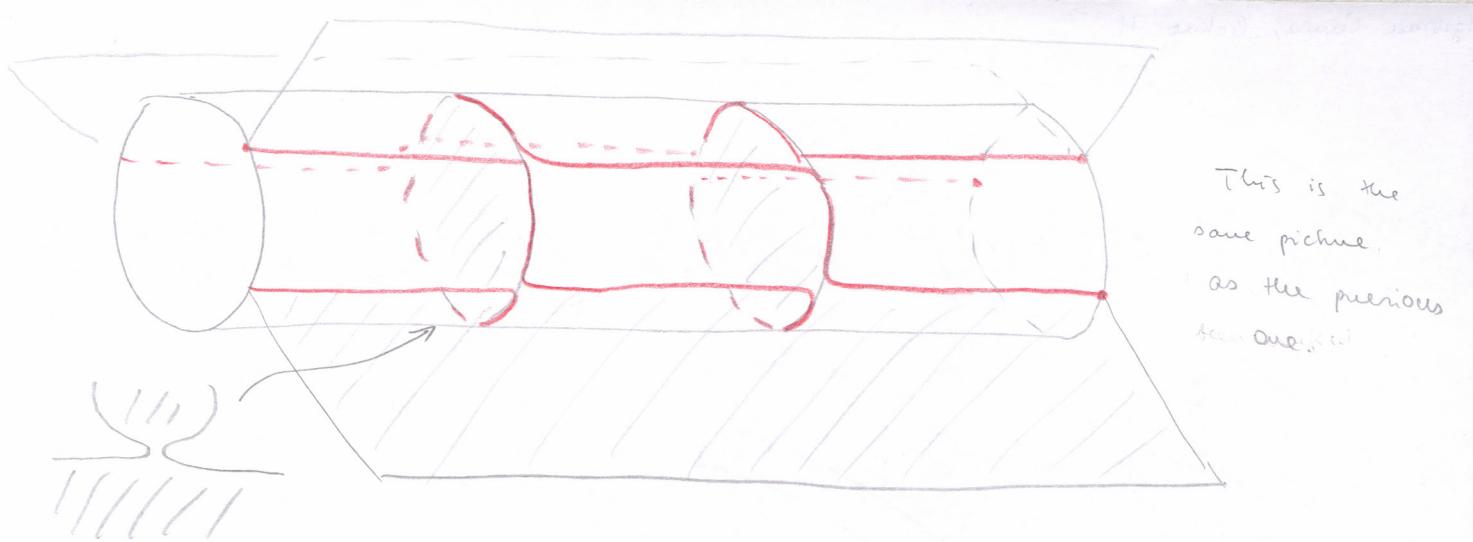
Blue: gluing points, around these:



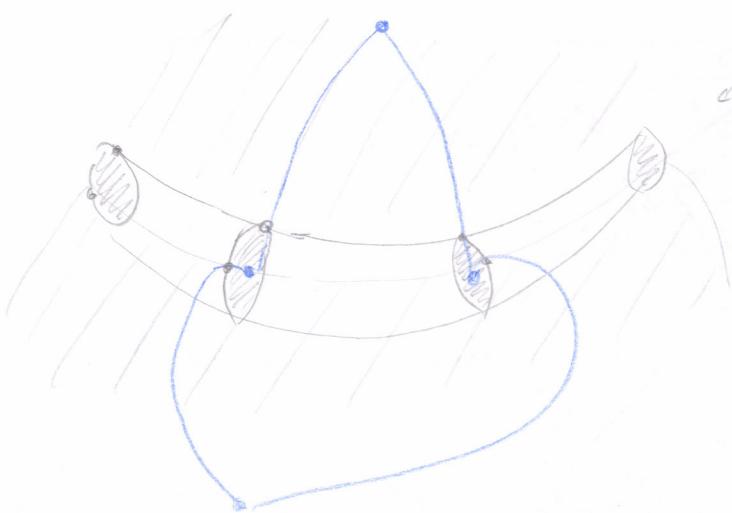
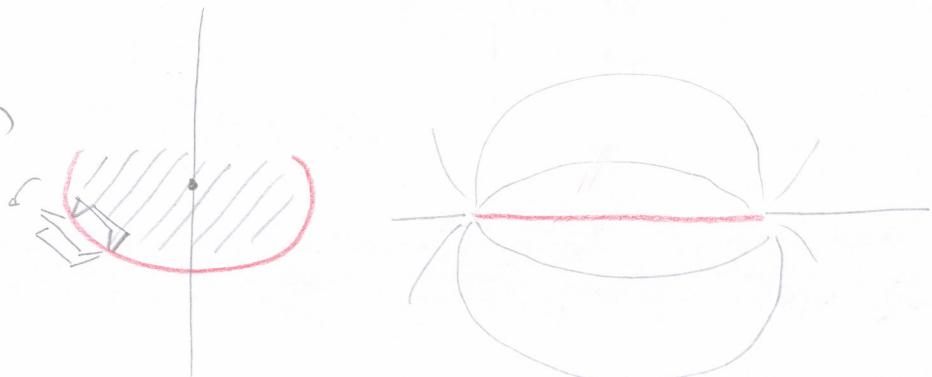
$a=3$



split these arcs  
in half, push one  
half downwards,  
the other one upwards  
to get the  
next picture



$$\tau(2,3) = \mathbb{O}_{(2,3)}$$



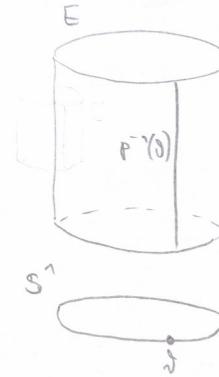
what the hell is this?

### Monodromy

$$F \hookrightarrow E$$

$$\begin{array}{c} \downarrow p \\ S^1 \end{array}$$

fibre bundle

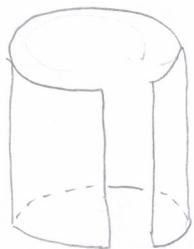


$$p^{-1}(s) \cong F$$

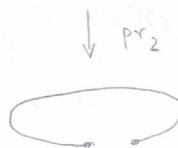
$$\text{In our case: } E = S^3 \setminus K$$

$F = \text{Seifert surface of } K$

Getting  $E$  open:



$$F \times [0,1]$$



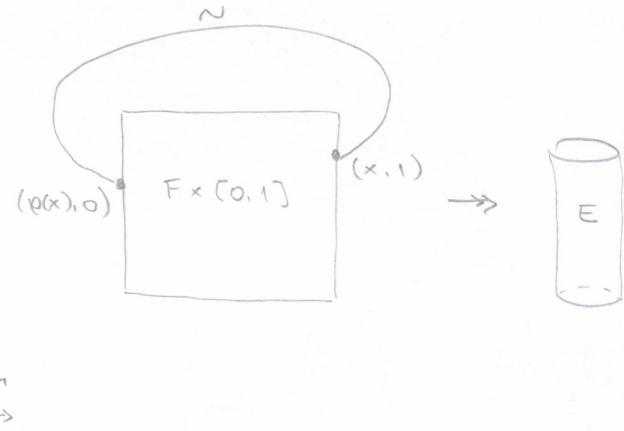
$$[0,1]$$

→ we get a bundle over  $[0,1]$

The fibre bundle  $p: E \rightarrow S^1$  is thus described by a map  $\varphi: F \rightarrow F$  called the monodromy of  $p$ .

(Note that  $\varphi$  need not be unique.)

$$\begin{array}{c} E \cong F \times [0,1] / (x,1) \sim (\varphi(x),0) \\ \downarrow G \\ \downarrow P \\ e^{2\pi i t} \end{array}$$



Constructing  $\varphi$ : let  $\theta_2 := 2\pi i t$ , this gives a vector field on  $S^1 \subset \mathbb{C}$

$\Theta :=$  a lift of  $\theta$  to  $S^3 \setminus K$

$$(dp)_x: T_x S^3 \rightarrow T_{p(x)} S^3$$

$$\underbrace{T_x F_{p(x)}}_{12} \oplus N_x F_{p(x)}$$

$p$  does not vary here,  
 $dp = 0$

$$\Theta_x = \begin{cases} \left( (dp)_x|_{N_x F_{p(x)}} \right)^{-1} (\partial_{p(x)}) & \forall x \notin K \\ 0 & \forall x \in K \end{cases}$$

(We won't need this explicit formula for  $\Theta$ )

Let  $\Phi_t: S^3 \rightarrow S^3$  be the flow of  $\Theta$ . (called the monodromy flow)

$\Phi_t$  maps  $F_{e^{2\pi i t}}$  to  $F_{e^{2\pi i (t+\ell)}}$  homeomorphically

$\varphi := \Phi_1|_{F_1}: F_1 \rightarrow F_1$  is the monodromy.

Example.  $K = T(a, b) = S^3 \cap \{y^a - x^b = 0\}$

Milnor showed that

$$p := \frac{y^a - x^b}{|y^a - x^b|}: S^3 \setminus K \rightarrow S^1 \text{ is a fibre bundle.}$$

$$\zeta_t := e^{2\pi it/a} \cdot \zeta, \quad \mathbb{C}^2 \longrightarrow \mathbb{C}^2$$

$$\Phi_t: \mathbb{C}^2 \longrightarrow \mathbb{C}^2$$

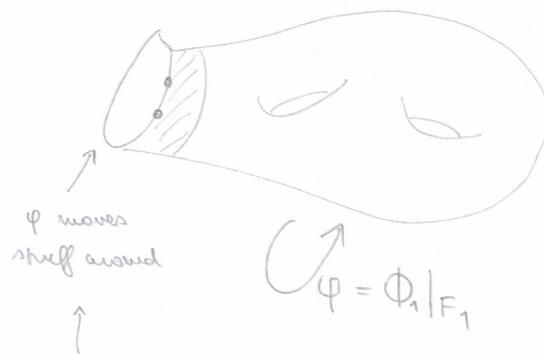
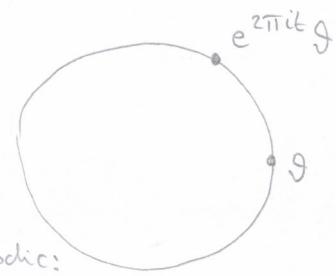
$$(x, y) \mapsto (\zeta_t^a x, \zeta_t^b y)$$

$$\text{Observe that } p(\zeta_t^a x, \zeta_t^b y) = \frac{\zeta_t^b}{|\zeta_t^b|} p(x, y) = e^{2\pi it} \underbrace{p(x, y)}_d$$

$\varphi = \Phi_1|_{F_1}$  is the monodromy. We obtain that this is periodic:

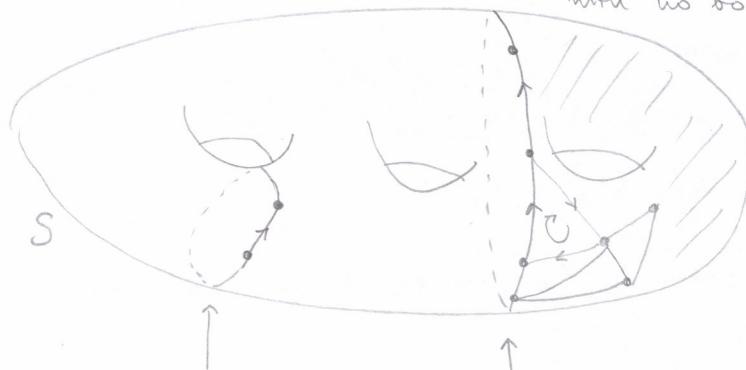
$$\Phi_1(x, y) = (\zeta_1^a x, \zeta_1^b y) = (e^{2\pi i/b} x, e^{2\pi i/a} y)$$

Thus the monodromy of a torus knot is periodic, up to an isotopy twisting the boundary.

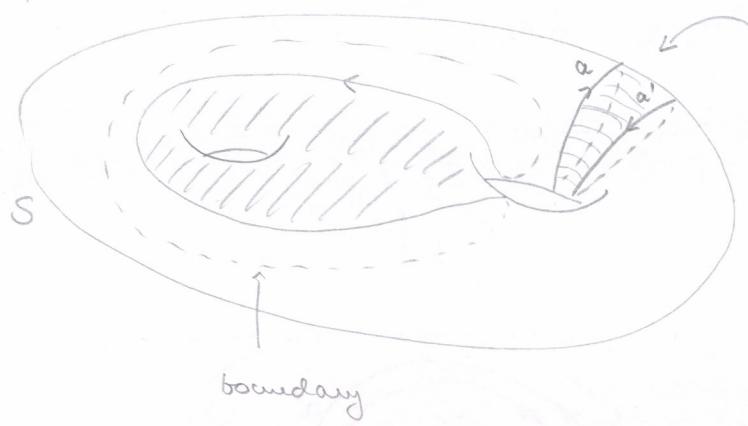


This can be remedied with an isotopy.

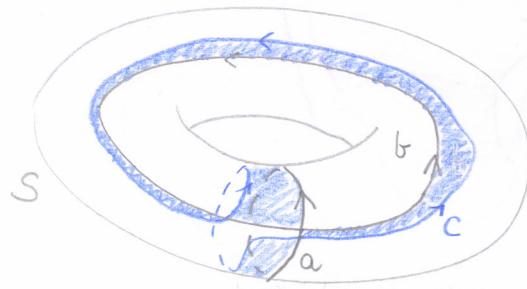
Recall:  $H_1(S; \mathbb{Z}) = \left\{ \begin{array}{l} \text{$\mathbb{Z}$-linear combinations of 1-dim oriented "pieces"} \\ \text{with no boundary} \end{array} \right\} / \left\{ \begin{array}{l} \text{edges of 2-dim "pieces"} \\ \text{with boundary} \end{array} \right\}$



Here we don't specify what "pieces" mean; this holds for several manifolds this.



$$\begin{aligned} 1 \cdot a + 1 \cdot a' &= 0 \in H_1(S; \mathbb{Z}) \\ \Rightarrow -a &\text{ is } a \text{ with inverse orientation} \end{aligned}$$

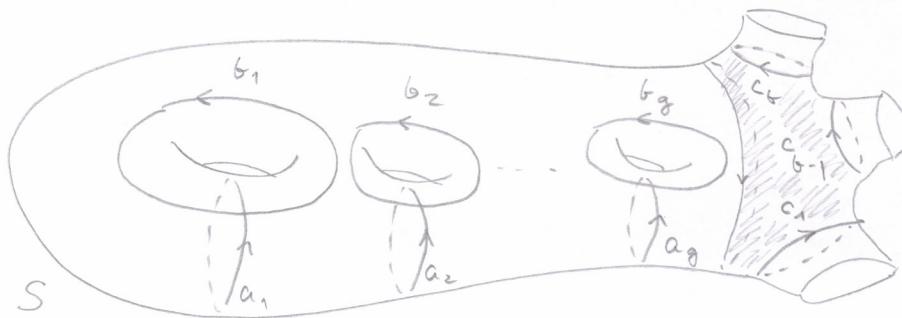


$$1 \cdot a + 1 \cdot b = 1 \cdot c$$



Fact:  $S$  compact connected oriented surface  $\rightarrow H_1(S; \mathbb{Z}) \cong \mathbb{Z}^{2g(S) + b(S) - 1}$   
is a  $\mathbb{Z}$ -module where  $g(S)$  = genus of  $S$ ,  
 $b(S)$  = # boundary components of  $S \geq 1$

(If  $b(S) = 0$  then  $H_1(S; \mathbb{Z}) \cong \mathbb{Z}^{2g(S)}$ )



We take all but one boundary curves since the sum of all boundary curves is a boundary, hence 0.

$$H_1(S; \mathbb{Z}) \cong (a_1, \dots, b_1, \dots, c_1, \dots, c_{b-1})$$

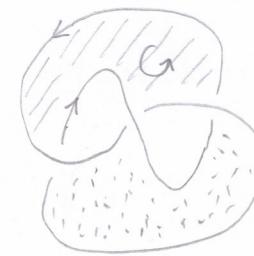
## Seifert surfaces

Let  $S \subseteq S^3$  be a Seifert surface.

The Seifert form of  $S$  is the bilinear

form  $(\cdot, \cdot): H_1(S, \mathbb{Z}) \times H_1(S, \mathbb{Z}) \rightarrow \mathbb{Z}$

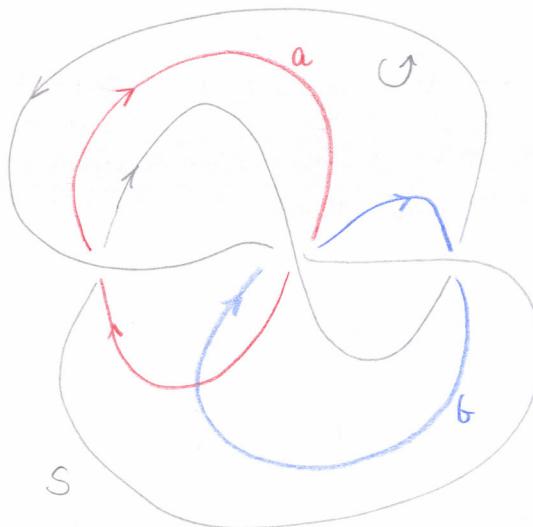
$$(a, b) \mapsto lk(a, b^+)$$



where  $b^+$  is the curve obtained from  $b$  by pushing  $b$  off into the positive normal direction to  $S$ .

(Note that since  $S \subseteq S^3$ , every point of  $S$  has a normal vector, and oriented)

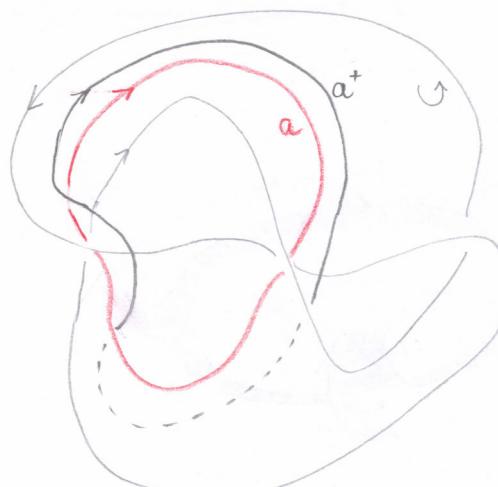
Ex.



$$H_1(S, \mathbb{Z}) \cong \langle a, b \rangle_{\mathbb{Z}} = \mathbb{Z}^2$$

Seifert matrix:

$$\begin{pmatrix} a^+ & b^+ \\ a & (-1 & 0 \\ b & 1 & -1) \end{pmatrix}$$



right-hand rule:



$\Rightarrow$  positively linked

YouTube: ddreibel pb23

3. h

Let  $S$  be a fibre surface,  $S^3 \setminus \partial S \rightarrow S^1$ ,  $\partial S = K$

$\varphi: S \rightarrow S$  monodromy map,  $\varphi = \varphi_1$ .

$\varphi_*: H_1(S, \mathbb{Z}) \rightarrow H_1(S, \mathbb{Z})$  induced map

$$[c] \mapsto [\varphi(c)]$$

Lemma.  $(v, w) = (\varphi_*(w), v)$   $\forall v, w \in H_1(S, \mathbb{Z})$

PROOF:  $(v, w) = lk(v, w^+) = lk(v, \varphi_{1/2}(w)) = lk(\varphi_{1/2}(v), \varphi_{1/2}(\varphi_{1/2}(w))) =$

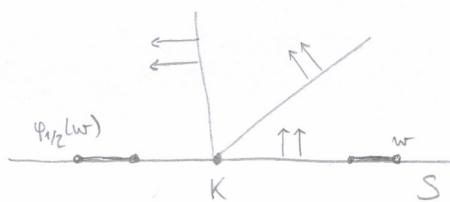
$\uparrow$   
 $\varphi_t$  is an isotopy, flowing does not change lk

$$\varphi_{1/2} \circ \varphi_{1/2} = \varphi_1 \quad (\text{flow}) \Rightarrow (v, w) = (lk(\varphi_{1/2}(v), \varphi_1(w)))$$

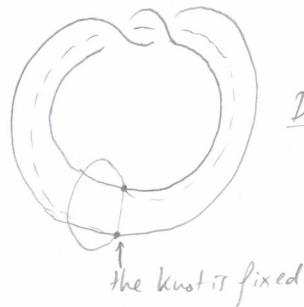
$$= lk(\varphi_1(v), \varphi_{1/2}(w)) \quad lk \text{ is symmetric}$$

$$= lk(\varphi_1(w), v^+)$$

$$= (\varphi_1(w), v)$$



example of why the monodromy map is not the identity:



it applies a Dehn twist.

the knot is fixed

example where monodromy map is Id: a disk



Choose a basis of  $H_1(S, \mathbb{Z}) \cong \mathbb{Z}^n$ . Write matrices  $A$  for  $(\cdot, \cdot)$ ,  $M$  for  $\varphi_*$ ,

$$(x, y) = x^T A y.$$

The Lemma says  $x^T A y = (M y)^T A x$ .

$$(M y)^T A x = \underbrace{y^T M^T A x}_{\text{scalar}} = x^T A^T M y$$

→ it is its own transpose

$$\Rightarrow A = A^T M \Rightarrow M = (A^T)^{-1} A$$

using that  $\det A = \pm 1$  (believe this for now)

Def. Let  $S$  be a Seifert surface.  $\boxed{\Delta_S(t)} := \det(tA^T - A)$  where  $A$  is the Seifert matrix. This is called the Alexander polynomial of  $S$ .

Thm.  $\Delta_S(t) = \pm t^{\frac{1}{2}\chi_K} \Delta_{S'}(t)$  if  $\partial S = \partial S' = K$ , i.e. if  $S$  and  $S'$  bound the same surface.

Def.  $\boxed{\Delta_K(t)} := \Delta_S(t)$  for some Seifert surface  $S$  of  $K$ .

Ex.  $\left. \begin{array}{l} \Delta_{T(2,3)}(t) = t^2 - t + 1 \\ \Delta_S(t) = 1 \end{array} \right\} \Rightarrow$  the trefoil knot is not trivial.

Thm. Let  $S$  be a fibre surface with monodromy  $\varphi: S \rightarrow S$ . Then

$$\Delta_K(t) = \pm \chi_{\varphi_*}(t)$$

where  $\chi_{\varphi_*}(t)$  denotes the characteristic polynomial of  $\varphi_*: H_1(S; \mathbb{Z}) \rightarrow H_1(S; \mathbb{Z})$ .

Proof:  $\Delta_K(t) = \det(tA^T - A)$  by def.

$$\begin{aligned} &= \underbrace{\det(A^T)}_{\pm 1} \det(t \cdot \text{Id} - \underbrace{(A^T)^{-1} A}_{M}) \\ &\quad \underbrace{\chi_M(t)}_{\chi_{\varphi_*}(t)} \end{aligned}$$

Rem.  $A - A^T$  is skew-symmetric; this turns out to be the intersection form for curves on  $S$ . It can be written as

$$\begin{pmatrix} 0 & 1 & & \\ 1 & 0 & & \\ & & \ddots & \\ & & & \ddots \end{pmatrix} \rightarrow \text{has } \det = \pm 1$$

$$\text{But } A - A^T = A^T M - A^T.$$

$$\rightarrow \det A = \pm 1$$

Cor.  $K$  fibred of genus  $g \Rightarrow \Delta_K(t) = \pm 1 \cdot t^{2g} + \dots$

Using this, one can show that certain knots are not fibred.

Polar coordinates

$$\rho: X = \mathbb{R}_{\geq 0} \times S^1 \longrightarrow \mathbb{R}^2 = \mathbb{C}$$

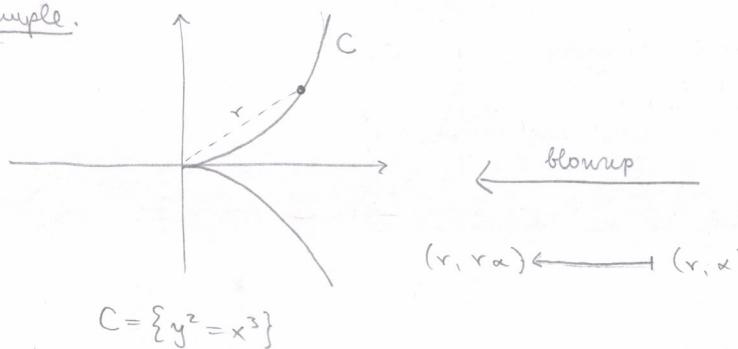
$$(r, \alpha) \longmapsto (r \cos \alpha, r \sin \alpha)$$



The only "bad" thing that happens here is that  $\rho(\{0\} \times S^1) = \{(0,0)\}$ , meaning that  $\rho$  is a diffeomorphism outside  $(0,0) \in \mathbb{R}^2$ .

$E = \{0\} \times S^1$  is called the exceptional divisor

$$\rho|_{X \setminus E}: X \setminus E \longrightarrow \mathbb{R}^2 \text{ diffeomorphism}$$

Example.

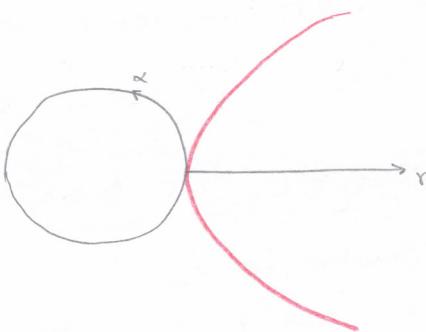
$$C = \{y^2 = x^3\}$$

$t \mapsto (t^2, t^3)$  parametrisation

$$r = \sqrt{t^4 + t^6} = t^2 \sqrt{1+t^2} \approx t^2 \text{ for } |t| \ll 1$$

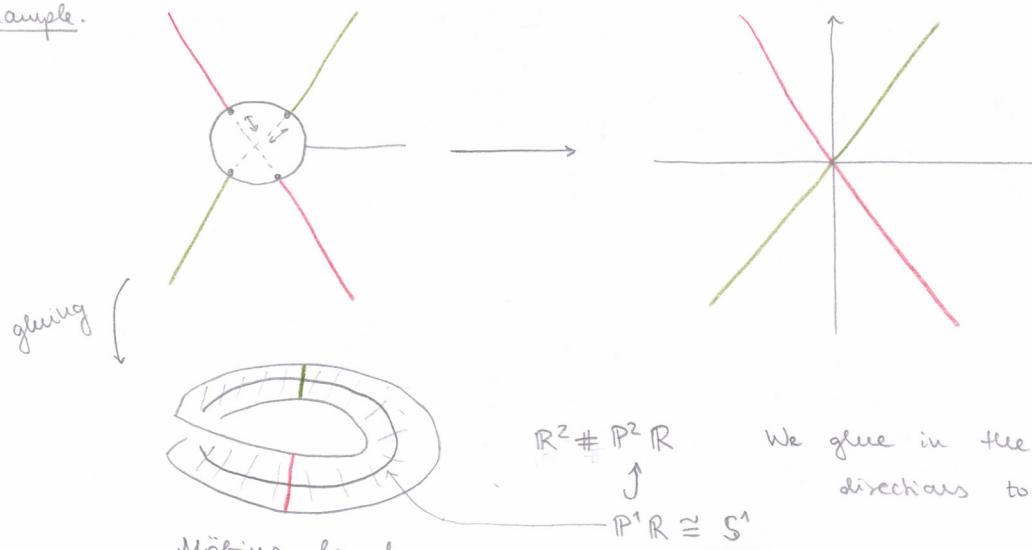
$$\tan \alpha = \frac{t^3}{t^2} = t$$

$$\tan \alpha = \alpha + \frac{1}{3} \alpha^3 + \dots \approx \alpha \text{ for } |\alpha| \ll 1$$



$$(r, \alpha) \approx (t^2, t)$$

This already exhibits the main property of blowups: they make singularities "less bad".

Example.

$$\mathbb{R}^2 \# \mathbb{P}^2 \mathbb{R}$$

$$\mathbb{P}^1 \mathbb{R} \cong S^1$$

We glue in the space of all directions to the origin.

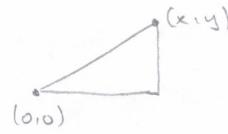
## Complex blowup

$$\mathbb{C}^2 \longrightarrow \mathbb{C}^2$$

$$(u, v) \longmapsto (u, uv) = (x, y)$$

in local coordinates

Globally:  $Y := \left\{ ((x, y), [\alpha : \beta]) \in \mathbb{C}^2 \times \mathbb{P}^1 \mathbb{C} \mid xy = \beta x \right\}$



$u = x$  is the radius

$v = \frac{y}{x}$  is the slope / angle

$$\sigma: Y \longrightarrow \mathbb{C}^2$$

$$((x, y), [\alpha : \beta]) \longmapsto (x, y)$$

think of this as the "algebraic form" of  $\frac{y}{x} = \frac{\beta}{\alpha}$

$$E := \{0, 0\} \times \mathbb{P}^1 \mathbb{C} \longmapsto (0, 0) \quad \text{exceptional divisor}$$

$\sigma|_{Y \setminus E}: Y \setminus E \longrightarrow \mathbb{C}^2 \setminus \{0, 0\}$  is an analytic isomorphism

$Y \cong \mathbb{C}^2 \# \mathbb{P}^2 \mathbb{C}$  as before (gluing in the directions at the origin)

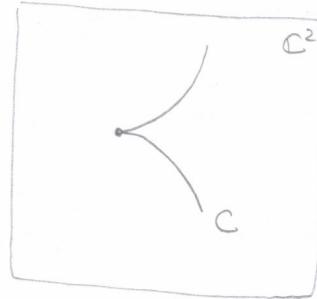
local coordinates for  $Y$ :

$$\varphi_1: \mathbb{C}^2 \longrightarrow Y \quad \begin{matrix} (u, v) \longmapsto ((u, uv), [1 : v]) \end{matrix} \quad \text{this is the sort of chart we live on } \mathbb{P}^1$$

$$\varphi_2: \mathbb{C}^2 \longrightarrow Y \quad \begin{matrix} (u, v) \longmapsto ((uv, v), [u : 1]) \end{matrix}$$



$$\sigma$$



$$C = V(f)$$

$$\sigma^{-1}(C) = \underbrace{\sigma^{-1}(0)}_{E \cong \mathbb{P}^1 \mathbb{C} \cong S^2} \sqcup C^*$$

$\varphi_1$

$$\mathbb{C}^2$$

$$\mathbb{C}^2$$

$$\begin{cases} f & \text{on } C \\ \text{other } \mathbb{C}^2 & \end{cases}$$

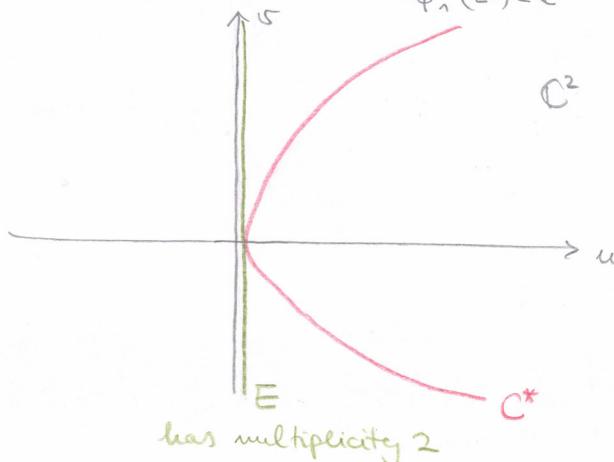
# Algebraic Curves, lecture 13

Let's look at  $\sigma^{-1}(C) = (f \circ \sigma)^{-1}(0)$  in the chart  $\varphi_1$ :

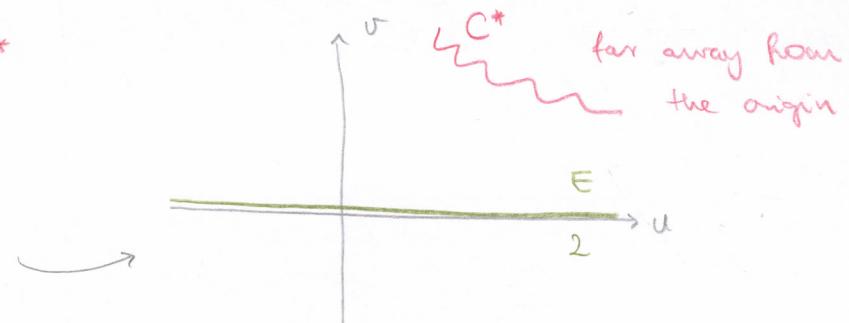
$$\{(u, v) \in \mathbb{C}^2 \mid f(u, uv) = f \circ \varphi_1(u, v) = 0\}$$

Example:  $f = y^2 - x^3$  (as usual)

$\varphi_1: 0 = f(u, uv) = u^2 v^2 - u^3 = u^2 (\underbrace{v^2 - u}_{} )$  this describes the union of 2 curves  
 $\varphi_1^{-1}(E) \subset \mathbb{C}^2$   $C^*$  strict transform of  $C$



$\mathbb{C}^2$  ← We are in  $\mathbb{C}^2$  again  
 $\Rightarrow$  we can blow up again!



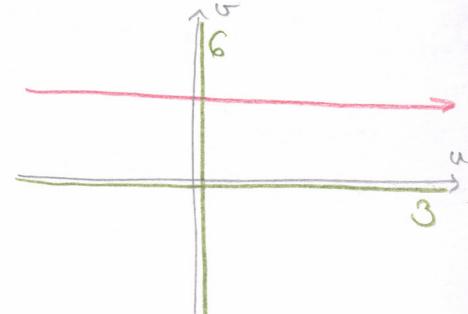
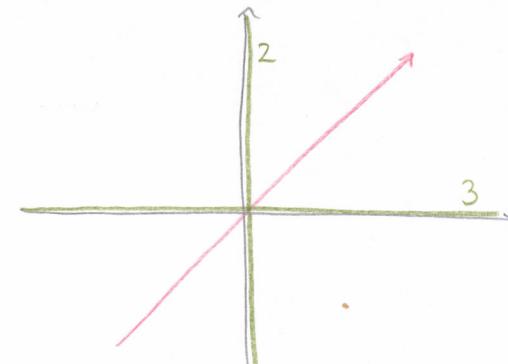
$\varphi_2: 0 = f(uv, v) = v^2 (1 - u^3 v)$

Now we blow up again:  
 $f_1(x, y) = x^2 (y^2 - x)$   
 $f_1(uv, v) = u^2 v^3 (v^2 - u)$

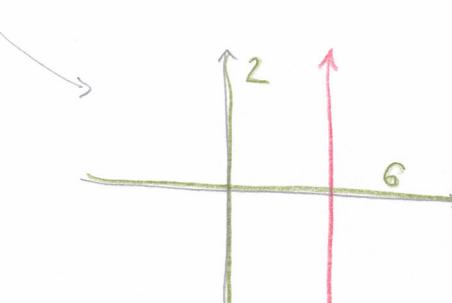
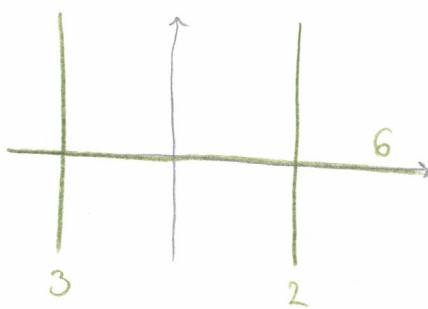


And again:

$$f_2(x, y) = x^2 y^3 (y - x)$$
  
 $f_2(u, uv) = u^6 v^3 (v - 1)$   
 $f_2(uv, v) = u^2 v^6 (1 - u)$



Summary:



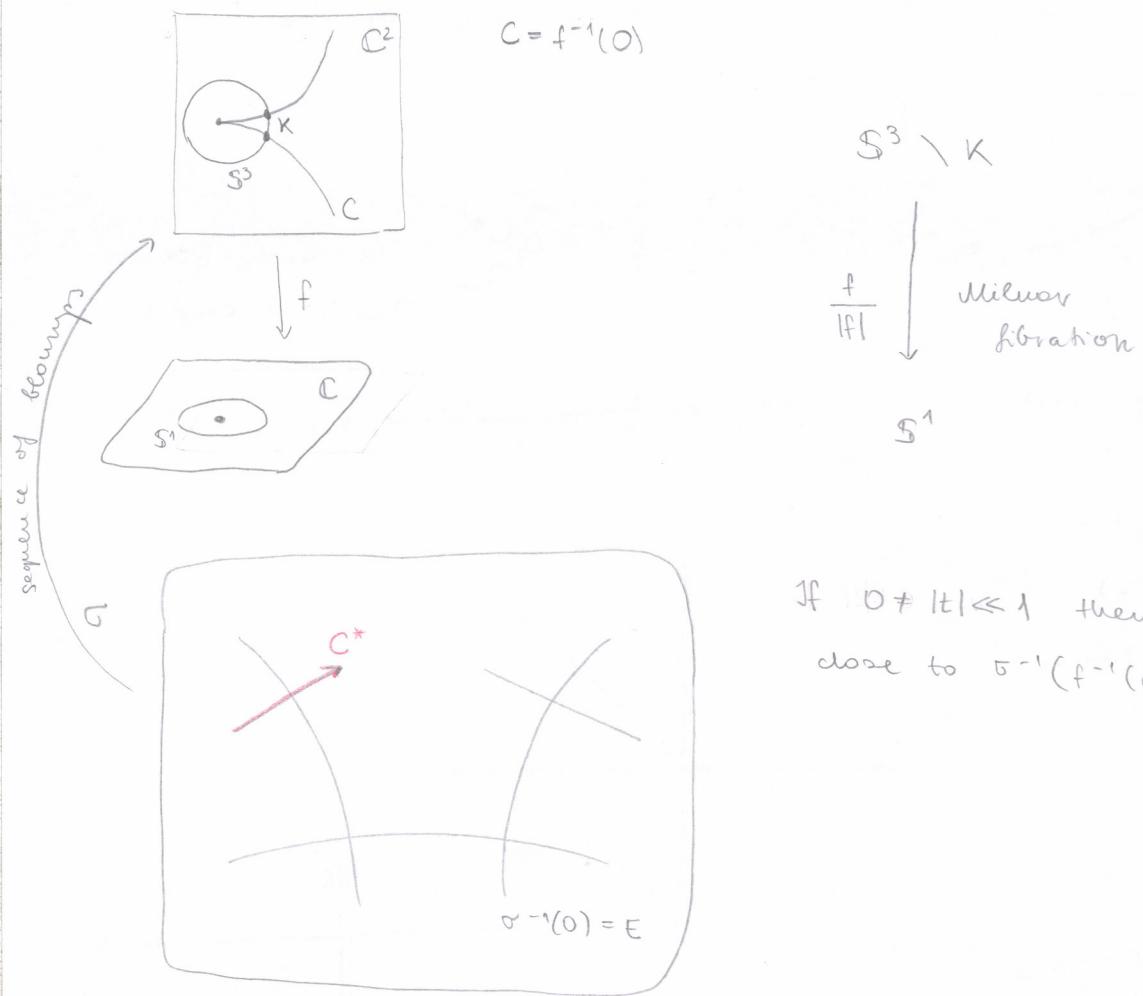
If we were to go on:  
 Quantitatively nothing changes  
 but the picture gets more complicated.

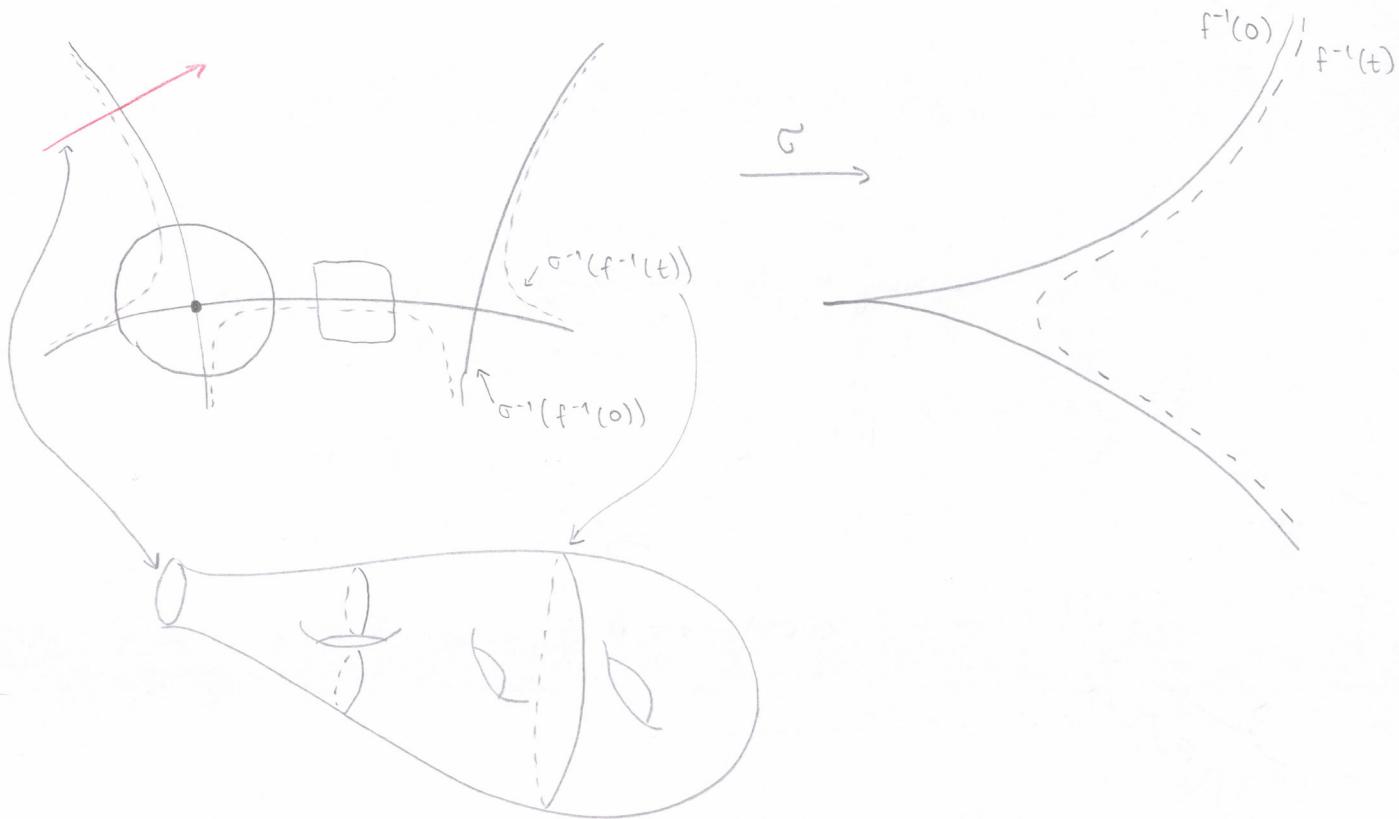


By resolving a singularity by a sequence of blowups, we mean arriving to the like we just did: every crossing is transversal, and everything is smooth. Newton's algorithm is somehow analogous to this, and it works for the same reason Newton's algorithm does.

Resolution of singularities: [Hauser et al.] is a nice survey of this proof. For fields of char 0: Hiraoka. In positive char, this is still open.

A'Campo 1975: understanding the monodromy of a singularity using blowups.

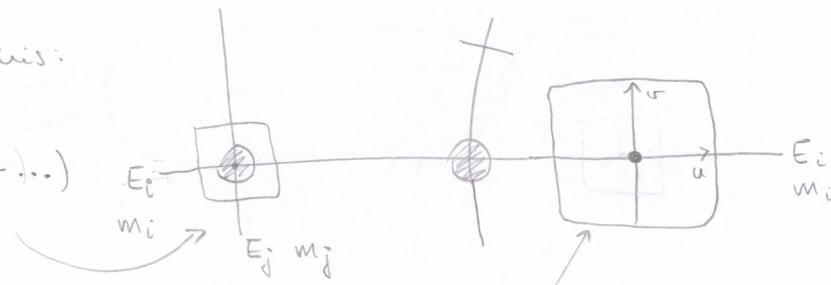




The blowup looks like this:

locally:

$$v^m = u^{m_i} v^{m_j} (1 - \dots) \quad (t =)$$



$$v^{m_i} (\dots) = 0 \quad E_i$$

$$v^{m_i} = t \quad \sigma^{-1}(f^{-1}(t)) = F_t$$

The monodromy is periodic outside the crossings



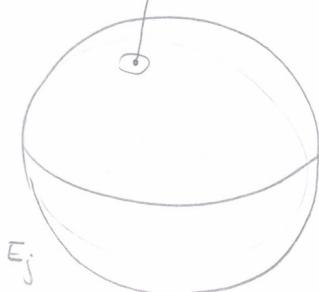
The fibre surface  $F$  of  $K$  decomposes into pieces

$$(subsurfaces) \quad F = \bigcup_{i \in I} F_i, \text{ one for each } E_i.$$

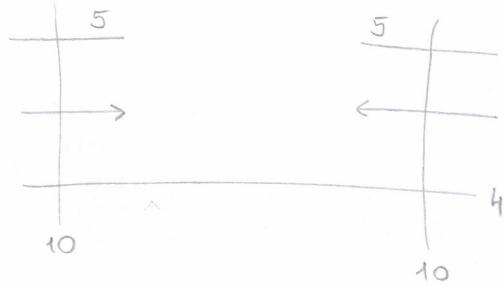
$F_i$  is an  $m_i$ -fold covering of  $E_i \setminus \bigcup_{j \neq i} E_j$

The monodromy  $\varphi|_{F_i}$  is the covering translation.

At each  $E_i \cap E_j$  we get  $\gcd(m_i, m_j)$  boundary components of  $F_i, F_j$ .



Example.



$$E_i \cong \mathbb{P}^1 \times \mathbb{C} \quad \text{Removing 2 pts: } \begin{array}{c} \text{---} \\ \circ \quad \circ \\ \text{---} \end{array} \quad \stackrel{=}{=} \quad \begin{array}{c} \text{cylinder} \end{array}$$

$$E_i \setminus \bigcup_{i \neq j} E_j \cong \mathbb{P}^1 \times \mathbb{C} \setminus \{\text{2 pts}\} \cong \begin{array}{c} \text{---} \\ \circ \quad \circ \\ \text{---} \end{array} \quad \text{pair of pants}$$

$$\chi(F_j) = 10 \quad \chi(\text{pair}) = -10 = 2 - 2g(F_j) - \#\pi_0(\partial F_j)$$

Result:

